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BDMS

Berkeley Database Management System
USER 'S MANUAL
(Version 2.2)

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# Berkeley Database Management System. <br> USER=S MANUAL <br> (Version 2.2) 

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## HARNING

This manual is pretimirary and in few cases describes features of BDMS that are not yet lully implemented. In particular,

1. double precision and bit string type data elemerts may be used procecurally, but not in the edltor and overy user languages,
2. DISPLAY and SUPFFESS commands are not yet lmplenented, *
3. the DUMP, CLEAN, and EALANCE utillties are not yet avallable.

BCMS had its origin ir the joint develoment of a farticle frysics Data System by tre Berkeley Particte Data Group and the Cattech Data Compllation Group. The diversity of databases that make up that system recessitated the use of a common database management system, while the complexity of the data demanded capabilities rot avallatie in any existing DBMS. The new system which was developed to meet this need has proven to be of general usefulness, and resporsibility for Its continued support and further development has been assumed by the Computer Scierce and Applied Mathematics Department at Lel.

As the system evolved, many people contributed ldeas and asslisted with programmirg. They inclute Tricia Coffeen, Paul chan, Geoffrey Fox, Marce Hutchirson, Tom Lasinski, Deane Merrili, Gili Ringland, Alan Rittenberg, Silvia Sorell, Paul Stevens, Ton Tripoe, Vicky White, and George Yost.

The users of BDRS also deserve thanks for contercing with the seemingly endess change to mich a develooing systen is sublect. Their comments, criticisi, and unfaling abillty to flust out even the most rectuse bugs have been and continue to te invaluable.

This work was suppertec by the U. S. Department of Energy fformerty U. S. Energy Research and Development Administratior and U. S. Atomic Energy Commissiorl, the National Science Foundation, and the National Bureau of Standards.

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### 1.1 INTRODUCTION

BOMS is a general-purpose database management and information retrieval system with a broad range of capabilities lor creating, maintaining, and eccessing computer databases. It frees programier and end user alike from concern with the physical storage of data, instead making it pcssible to deal with data at a logical level. In other words, it allows one to deal with information rather than data.

Many applications, which previously would have reduired extersive special program develcpment, can be handled by the syster with little, if any, extra software. Its capabilities include -

1. A natural anc easy-to-use database deflnition larguage.
2. A powerful editor that operates directly on the catabase. Data are enterec in a sinple, natural way, and mary short-cuts are avallable to expedite the entry of large amounts of data. Modifications to existing cata are immedztely effective.
3. Extensive retrieval facilities lncluding controlled detabase Inversior, Bcolean and relational operators, rested parentheses in search expressions, truncated and range search, and reference to previousiy retrieved sets.
4. A stancard listing format that makes the structure of the data reacilv apparent. When working interactivelv, retrieved records may be listed at the terminal or orinted off-line. They alsc ma, be dumped in a format suitable for data exchange and database reloading.
5. A common command and data language for both batct and interactive lse.
6. Exits to user-supplled routines at several places in the system to allow input data valldation, data transformation on input, outcut, indexing, searching, and autoratic data element generation.
7. Utilities for lading and dumping databases and databese space maintenance.

A designer wishing to base a speciel-puroose system on BOMS will find it relatively easy to use fost those modules neeced. All system capabilities are available via foftran subroutine calls. For example-

1. Soecial software might be writter to bulld a catabase from data lnitielly stored in some special format, e.c. fixed fielc card lrages. This could be done by writirg a orogram which reads date in the special format and calls appropriate subroutines to lcad it directly into the database, byonssirg the standard data input language. In fact, suct a routine could be coded as a replacement for the standarc input routine with the system behaving normally in all other
respects.
2. The system could be modified through the provisicn of a reptacement for the standard data output routine so trat retrievec cata would be listed in a format better sulted to its enc use.
3. A very arbitious system deslgner might choose to replace the entire user interface with one providing extersive prompting and data valldation tailored to his specific application.

In all of these cases, the designer/programmer is freed from concern with physical storage mechanlsis for his data. Insteac, he can use existing software to access and modify the date, referring to it by name.

BOMS ls righly nodular and coded for the most part in machine-indepencent FGRTFAN IV. The operating system irterfece and machine-dependent code are isolated in a few modules so thet the system is easily transportable.

## 1.2 dataease structure

A BDMS database is structured into records, the units in which deta pass between the system and jlisc storage. Normally, a record wlll have some significance to the user, e.g. a record in cibliographic database would be a descriotion of a single document, but this is not always necessary or desirable.

The individual data items within a record are called sfta elements: they are the smallest inits of data with any meanirg tc the system, although an individual data element might have some interral structure known to an apollcation program. A data element has a uniaue name and is rormally referenced by name (or a synonym). There ls essentially no limit on the number of data elements that may be ceflned for a database.

A blerarchical siructure may be lmposed on the records wher a database is cefinec. This means that some data elements are dectared to be sybordinate to other, carent, data elements. Those deta elements for which no parent is dectared are called recgrd-leyel deta elements; it is often useful to consider the record liselito te thelr parent. There is essentialy no limit on the number of levels whlch may be defined in the record structure.

Within a given record, eacr record-level data element may occur once, several times, or not al all. Likewise, each occurrence of a data element at any lower level in the nierarchy may nave linked to it one, several, or no occurrences of each of its sutcrdinate data elements. Jhere is essentialiy no limit on the number of times eny data element may occur in a single record. Furthermore, there is no storage overhead associated with data elements that do not occur at alf in elther the recorc or a particular occurrence of thelr. parert.

Data elements are clessified irto six types according to the yalues they can assume - character or blt string, integer or real isingle or double precisiorl vector, and pure node. Character or bll strings may be of any length with no limit beyond that imposed by run-time memery restrictions. Intecer or real (floating ooint) date elements may be scalars (single runbers) or arbitrary length vectors (i.e. an ordered set of numbers, which are the components of the vector). Real deta elements may be single or double precision. Pure node cata elemerts carry no value; they may be used to link together subcrdinate data elements in the record hierarchy or as flags.

Ary data element, regaroless of type, may serve as a node in the hierarchical recorc stricture. In general, if ore of a group of related data elements may occur only once in each occurrence of the group, it should be made the parent of the rest of the data elements. However, if all of tre data elements in such a group mey occur multiolv, they shculd all be linked to a cure node parent.

Ary data elemart may be declared to be a record-key. The systeft will then meintalr an irdex for that data element to allow efficlent retrieval. In ar index, key values have a fixed ifngth trat is declared in the database defirition. Data efement values are
truncated or padded as necessary when they are put irto an index.
The system assigrs a record ID to each record as it is created. This guarantees that each record has a unique identilier by which it can be selectively retrieved even if none of its data elements is defined to be a key so that no indices are mainteined.

Physically, a database is divided into a data file, which contalns the database definition and data records, a directory file used to access recoros in the data file, and an inversion flle comprising the indices (if any exist).

### 1.3 EXAMPLE DATABASE

Throughout this manual, we shall use for illustration a hypothetical databise whose records are summaries of some kind of experiment reports ard wrose structure may be diagranmed as follows -
( FECORD )


The cata elements rave the following meanings -
AN is a document accession rumber.
TABLE is the rame of a collection of data. It links togetrer definitions of the $X$ and $Y$ variables, $X N$ and $Y N$, possible commerts on the tabie, $C T$, and the actual data polnts in the table, cATUM. DATUM is a pure node data element that inks $X$ and $Y$ values and possible comments on the datum, $C D$. $O Y$ is the error on $Y$. In a given record, TAELE night occur multioiy; in a given taEle, DATUM and CT might occur maltiply; in a civen DATUM, CD mignt occur multioly. The rest of the data elements; $X N, Y N, X, Y$, and $C Y$ would normally occur only once withir a given occurrence of their parent deta element.

AUTHORS is a pure node that links one or more autrors, $A$, with their institutioral affillation(s), I. In general, all of these data elements might cccur multiply. Note how PRINCIPAL, a pure node type data element, is used to llag a particular author as the princical author. Normally, only one author would be so flegged.
$R$ is the cocument reference (e.g. report number) and $T$ and $\bar{D}$ are its type and cate.

A diagram of this type only shows the hierarchical relatiorships among cata elements. Ir general, any of the data elefments may occur several times or not at all in any oarticular recorc. trus, the pictorial representation of a real record would require the use of a thiro dimension to cisplay the multiple occurrences.
1.4 SYSTEM ORGANIZATICN

From the user*s point of view, gDMS is divided into several functional modules.

The gatabsse definition comolier is a separate frogram used to create a new database before any data are loaded into it. The same editor used for data input and modiflcation is usec to construct a database definiton trat is converted by the comoiler into tables trat will drive the rest of tre system when the new database is used. The database definitior language anc use of the compliter are the sublect of Chapter 2.

The sxecutive is the overall control orogram. Some executive commands are exected immeditely while others cause control to be passed to other system modules, which will then interoret additloral user input. The exective functions are discussec in chapter 3.

The editor provices facllitles for creating aro foditying deta records. It is controllec by a set of editor commands anc understands the system*s external data larguage. The use of the editor ls described in cetall ir chapter 4.

The guery subsystem, descriced in Chapter 5, interoretsuser queries and selectively retrieves data records for subsecuent display or modification.

The utlilities, described in Gmater 6 , are stand-alone prograns for database maintenance operations such as efficient loadirg and dumping and removal of tie dead space resulting from updete activity.

### 1.5 USER EXITS

As mentiored in the introduction, exits to user routines ere orovided at several pleces in tre system, allowing lt tc be talloned to specific databsses and applicatlons.

The subroutines that may be supplied by the user ere descrited oriefly below. Details of their use and calling parafeters ray be found in the Frogranmers Manual.

IPROC is passed a data element valuebefore it is stored in the record. IPFOC may modify the value or check 14 for validity: I e value is found to be in error, IPROC cen instruct the calling routine not to store it, ard/of outout an error message.

OPROC is passed a date element value before it is cutcute it mey modify the valve, check lf for validity, or suppress its output altogether. Normally, IPROC and OPROC would be used as a pair to transform between the internal and exterrel forms of a data element.

KPROC is passed tre velue of a key data element before it is storeo in an index lor deteted from an index wren restcring a modiflec record or adding a recordl. It mav rodify the value or check for validity. It can suppress tre storage of a key ir an index.

KMAP is passed tre value of a CHAR tyoe key data element before $1 t$ is stored in or deleted from an index. It rive perform an arbitrary character code mapping in order to enlorce a desired ccilating sequerce for values in the index.

QPROC is passed the value ol a key data element before an index search. Normally it would transform the value in the seme way as KPfOC, Iftmight consist of no more trar a call to KPRCC) but it is crovided to allow the use of cifferent forms of a key value during data input and retrleval.

SPROC is called before a record is stored in the cetabase. It ray perform an Intre-record data integrity check, ard generete or modify data element values. It can orevent the storage of the record il errors are detected.

FPROC is called lust after a record is fetched from tre database. Its prinary use is to rematerlelize any virtuel data element occurrences that the detabase designer wlshes to make visable to the user.

The IPROC, CPROC, KPRCC, and GPROC routines are callec only if the. database definition instructs the system to do so. kNAP, SPROC and FPROC are alweys callec but, of course, may be do-notting routines.

### 2.1 INTRODUCTION

To set up a BCMS catabase, one must define the nature of the cata, e.9. names and tyoes of data elements, their hierarchical relationships, whet indlces are to be maintalned. etc. The system must be informed of any user-supplledroutines for input processing and validatior, outout, key or auery processing.

All of thls information ls coced in a database definition lanouage which is processec by the database definition compiler. The outout of the complier is a binary lile definltion table (FDT) whict becomes the zero*th record cf tre newly-defined database anc controls the operetion of the system whlle that database ls being accessed.

Plans exlst for extending the complier to provide fecillties for modifying the definitior of ar existing databese eog. adding or deleting data elements and synonys.

### 2.2 DATABASE CEFINITION LANGUAGE

The catabase definitior language of Boms has the seme syntax as the data lancuage usec by the editor (c.f. Chapter 4). That is, it consists of a series of statements of the form

```
<attribute label> = <value>;
```

By way of illustration, a fige definition for the example database described in chapter 1 would be coded as follows, assumirg that it is desired to index the records by the document accessior rumber, an, the individual euthors, $A$, and the type of datareported, $Y N=$

```
FILE=EXPT-REFOTTS;
DE.=AN; SYN=ACCESSION-NO; TY=INTEGER; KEY; LENGTH=1;
DE.=TABLE; TY=CHAR;
DE.=DATUN; PAR=TAELE; TY=NODE:
DE.=X; SYN=X-VALLE; FAR=DATUN; TY=REAL;
DE*=Y; SYN=Y-VALUE; FAR=DATUN; TY=REAL;
DE* = DY; SYN=Y-ERROR; PAR=Y; TY=REAL;
DE.=CD; SYN=COMMENT-CATUM; PAR=DATUM; TY=CHAR;
DE.=XN; SYN=X-NANE; FAR=TARLE; TY=CHAR;
DE.=YN; SYN=Y-NAME; FAR=TARLE; TY=CHAR; KEY; LENGTH=1;
DE.=CT; SYN=COMMENT-TAELE; PAR=TABLE; TY=CHAR;
DE.=AUTHORS; TY=NODE;
DE.=A; SYN=AUTHCR-NAME; PAF=AUTHORS; TY=CHAR; KEY; LENETH=3;
DE.=I; SYN=INSTITUTICN; PAR=AUTHORS; TY=CHAR;
DE* =PRINCIFAL; PAR=A; TY=NODE;
DE.=R; SYN=REFERENCE; TY=CHAK;
DE=T; SYN=TITLE; PAY=F;; TY=CHAR;
DE*=O; SYN=DATE; FAR=R; TY=CHAR;
```

The meaning and allowed values of the databese defintion attributes are -

## FILEasfile names:

This is not yet used by the system.

## DEscpreferred data element names

Thls is the rame that hll be used when the data element is output. Names must be unicue within a database. They may not contaln embeded blanks, equal signs, or semicclons. The slze of tre FDT wlll be minimized and the editor's processing of input to the defined databese will be silghtiy more efficlert if the data element names all nave similar lengths li.e. all it into the same number of computer words).

## SYN:esynonym for data element names:

Synorvms may be usec interchangeably with the oreferred data element name for input and in aueries. They must satisfy the same rules of constriction. Again, it is preferable for all synoryms to have sillilar lergths.

## PARz<parent data efement preferred name>;

This specifles the parent to which a data element will be linked in the hierarchical necord structure. The parent data element must heve been deflined previousiy. PAR is omitted for record-level data elements.

## TYasdata element type>;

Allowed values are -

| INTEGER | irteger (vector) |
| :--- | :--- |
| REAL | real (vector) |
| DOUELE | dcuble precision real (vector) |
| CHAR | craracter string |
| BIT | bitstring |
| NODE | pore node |

## VIRTUAL:

If this follows a $D E$ specification, the data element will be discarded when a record is stored in the database. It ray be input, however, and may be usec as a KEY. It will stlli exist in the record when the SPROC routine is called. If it is desirec tret a VIPTUAL data element be rematerialized when a record is fetched from the datacase, this may be dore by the FPROC routine.

KEY:
If this follows a $D E$ specification, an index will be maintained for the data element.

LENGTHaskey ength in words?
This may fcilon a KEY attribute. If absent, trekey lencth defauts to 1 word.

## IPROCE<format of DE value>?

If this follows a DE specification, the IPROC routine will be called by the ecitor wren the data element is input. Allowed format specifications are -

EXTERNAL The DE value is passed to and is returned by IFFOC in external format, l.e. unpacked ctaracter string.

INTERNAL The $0 E$ value is passed to and retornec by Iffoc ir internal format, l.e. internal birary rumeric representatior, packed character strirg, or bit strinc.

CONVERT The DE value is passed to IPROC in external format and returned by IPROC in internal feriat.

## OPROC=<format of $D E$ values:

If this follows a $D E$ specification, the OPROC routine will be called by the list routine before the deta element is outcut. Allowed format soeciflcations are the same as for IPROC except thet the CONVERT has the opposite meaning i.e.

CONVEFT The DE value is passed to OPROC in internat format and returned by OPROC in external format.

KPROC=<format of DE value>?
If this follows a oE specification, the KPROC routine will be called orior to storing a DE value in an index. Allowed format specifications are the same as for IPROC.

QPROCEeformat of DE values?
If this follows a DE specification, the QPROC routine will be called by the query interpreter when it encounters the $D E$ in a auery. Allowed format specifications are the same as for IPRCC.

IPROC; OPROC; KFRCC, and QPROC are not valid attributes for a oure node type $D E$.

### 2.3 COMPILER COMMANDS

The commands necognized by the compiler are described ir this section. Only the first four letters of a command are necessary for recognition.
define
This command invokes the editor, which initially will be in apoend mode, so that a new database definition may be input. This ereates a database definition record in core which has the following hierarchical structire -


VIRTUAL and KEY sre pure node tyoe and LENGTH is integer type. All other data elements are Cliff tyoe.

As this record is beirg input, errors or omissions may be remedied using the edt mode commands *i, ${ }^{*} R$, ${ }^{*} S$, and $* 0$ just as if it were a data record. It may be listec at any time with the ecitor command * L. When the definition is complete, the editor must teftermineted by a ** command.

## conpile

This commend iritlates compliation of the catabase definition created by a preceding DEFINE command. The complifer resporios with COMPILING DEFINITION
followed by either
COMPILATION SUCCESSFUL
or a list of errors detected and
COMPILATION JERMINATED

## MODIFY

This command invokes the editor to allow correctior of the errors detected in a prececing compilation. When the changes rave been made,

```
the editor must be terminated by a ** command. It differs from a
DEFINE command in that DEFINE wIll comoletelv erese an existing
definition record whlle MoDIFY allows changes to an exlsting
definition record.
```


## LINE, <n>

This sets the input/outout line length to en> characters. The defautt is 80 characters.

STOP
This terminates execution of the complier. If the definition has been complifed successifully an enpty datatase wlll exist ready for the addition of data records.

The audit facility is alway turned on when the database definition comoller is runt all input to the system from the Infut file the terminal when running interactlvely) is echoed on tre Audit lile. This provides a record of the fefinition that may be edited and used as batch-mode inout to the comoller if the definitior needs to be changed later. This copy of the definition might alsc be saved so that it.could be used when relcading a dumped copy of the detebase. e.9. Following transmittal to ancther site.

### 2.4 COMPILER ERROR MESSAGES

The error messages that may be gener ated by the database defirition compller are summarized in this section. The prefix **FEFAOR***, which is common to ill messages, has been omitted.

## Illegal recurrence cf reccro level atjribute

One of the ettribttes FILE, or DE has recurred without following period. It is igrored.

INVALID COMMAND
A command has beer ircut whose first four letters co not match any of those in Section 2.3.

INVALID FARENT SPECIFIED FOR DE.n
The parent soecilled for DE.n has not been deflned orlor to defining $D E \cdot n$ or is misspelled.

INVALID TYPE SPECIFIED FOf DE. $n$
Velid types are INTEGER, REAL, DOUBLE, CHAR, BIT, or NOCE.
INVALID IPROC TYPE SPECIFIEDFOR DE, $n$
INVALID OPROC TYPE SPECIFIEO FOR DEA
INVALID KPROC TYPE SPECIFIEE FOR DE, $n$
INVALID QPROC TYPE SPECIFIED FOR DE. $n$
Valio tyoes ere EXTERNAL, INTEFNAL, or CONVERT.

## no data elements oefined

The databise definition contalns no $D E$ attribute specification statemen is at all.

NO TYPE SPEGIfIED FGR DE on
Every data element mist have a tyde soecification.
PROC SPECIFIED fOF PUFE NCDE DE.n
Since a pure node type cata element carries no value, it is meaningless for a processor to operate on it.

WORK SPACE EXCEEDED -- CCWMAND SKIPPED
The database deflnition is too large to complle in the availatle work space. The compller must be recomplifed with a larger work space before proceeding.

## HORK SPACE EXCEEDED -- RUN ABORTED

There was not even enough work soace to initialize the system. This is an unlikely cccurrence indicative of syster mafunction.

### 3.1 INTRODUCTION

The EDMS executive is a self-contalned facillity that allows a user to access and maintein an exlsting database. Commands are provided to add new recoros to the databese, to search for recorcs satisfing stated criterla, and to subsequentiy display, orint, dump, modify, and delete such records. In additlon, several commands allow user control over certain escects of the system's operation such as input/outout, line length, and whether to mainteln an audit trail.

An overview of the executive wll be presented in this chapter, and the use of most commands will be descrited in detail. Two areas are sufficientiy irvolved to warrant separate treatment in the following two chapters - the editor and ouery language.

### 3.2 EXECUTIVE COMMAA DS

All executive commands take the form of an English-languege verb followed optlonally by a comma and an integer ldentifviro the record to be acted upon. The commands are terminated by a blank and hence may not contaln embedded blanks.

Whenever the executive is waiting for a user command, 14 issues the message

ENTER COMMAND
The commano verbs recognized by the executive are -

| ADD | Add new records to the catabase. |
| :---: | :---: |
| SET | Retrieve and make current a set of reat Make a previously created set curren |
| PURGE | Purge all previcusiy created sets or the last set created. |
| LIST | List at the terminal the current set or a selectedrecord. |
| PRINT | Print off-line the current set or a selected record. |
| DUMP | Bump the current set or a selected record in external format. |
| OISPLAY | Set default list of data elements to be listed, orinted, or dumped. |
| SUPPRESS | Set default Ilst of those data elements not to te 1 isted, printed, or dumped. |
| modify | Modify a record selectec from the current set. |
| delete | Delete irom the database a record selected from the current set. |
| LINE | Set irput/output line tength. |
| AUDIT | Turn on audit facility defautil. |
| NOAUOIT | Turn off aucit facility. |
| STOP | Terminate program execution. |

In somewhat more detail, these commands function as follows. ADD

This command irforms the execulive that one or more new records are to be added to the dat abese. An empty record will te created and assigned the next avallable recorcid. Control is then transferred to the editor so that the user mey enter data in the reccrdis). (see Chapter 4 for instructions on the use of the editor.l When the user exits from the ecitor (with a ** or ${ }^{*}$ C editor command), the number of records added to the database is reported ir the form
<n> RECORD(S) ADDED

## FIND

Tris command lrvokes the query subsystem and is lollowed by a condition terminated by a 7 \# delimiter. The database will be searched for records setistying the condition and if any are found, they form a new current set on which subsequent LIST, PRINT, CUHF, MOOIFY, or DELETE commands will act. The set identifier and number of records in

The set are reportec to tre user in the form
<n) RECORD(S) IA SET <S>

If $\langle n\rangle=0$, no set is createc and $\langle s\rangle$ is omitted. The format of queries will be covered ir detaitin Cheoter 5.

SET.\&set number>
This command makes a previously-created <set number> the current set on which subseavert LIST, PRINT, DUMP, MODIFY, and CELETE cormands will act.

## PURGE,<set numbers

This command purges areviousiy-created eset numbers, frefing the disc and work scace it uses. If <set number> is omitted, all existing sets are purged. (At present, only the lastecreatec set. or all sets. may be purged with this ccmmand.l

LIST,<n>.<data ellement nawe $1>0 . . .<d a t a$ ellement nane is

This command lists the neth record in the current set on the system file LOG (which is the termina: when running interactivelyl. If the number <n> is omitted, all records in the set will be listed.

Each data element occurrence will begin on a new lire and will be numbered if several occurrences of that data element are linked to the same parent occurrence, or if it is a record level data element whicr occurs more trar orce Subordinate data elements will be ilsteo following the parent occurrence to whlch they are lirked and indented according to their level in the record structure.

If the list of data element names is omitted, the last DISFLAY or SUPPRESS command issuec controls which data elements are displayed. If no DISPLAY or SUPPRESS command has been issued, all data etemerts are displayed.

If a list of deta elemert names separated by conmas follows the command verb (and record number en>, If presentl, orly those data elements in the list will be displayed. The entire conmend, including the list of cata element na ${ }^{\prime \prime}$ es, must not contain embedded blanks.

PRINT, <n>, <data element name $1 \geqslant 0.0<d a t a$ element name is

This is identical to LIST, except that the recordis) is ilisied on the system file FRINT, which can be disoosed for printing ofiline when the job is concluded.

DUMP,<n>, $<d a t a$ element nawe $1>0$...edata element neme is

This commano is similar to pRINT, except that the recoro(s) is dumped on the system file DUMF in a format readable by the editor. This file may be used subsequentiy to load anotrer detebase.

This command sets the datz elements to be listed, orlnted, or dumped by LIST, FRINT, or DUMF commands that are not lollowed by a list of data element names. This default is effective intil arctrer DISPLAY or SUPFRESS command ls issued. The data element names in the list are seoratec by commas. The entire command must not contain embecded blanks.

## SUPPRESS,<data element name i>g....<data element nane is

This command hes a function similar to the BISPLAY commard. The difference is that a SUPPRESS command specifles which data elenerts are not to be listed, printec, or dumped by LIST, FRINT, or DUMP commands that are not followed by a ist of data element nemes. This default is effective until another DISPLAY or SUFPfESS command is issued.

If no list of data element names follows a SUPPRESS commenc, the default is reset to all data elements.

## MCBIFY, <n>

This command ceuses the ens"th record of the current set to be fetched from the databise intc core and then passes control to the editor. The record can be modifled as desired using editor commands. When all editing is comolete, control is returnec to tre execullve, which restores the record in the database and updetes all indices to reflect the cranges in the record. Successful completion of this oderation generates the ressage

RECORO MODIFIED

## DELETE,<n>

This causes the en>*th record of the current set to be removed from the database. All indices will be updated to reflect thls action. Successful completicn gererates the message

## RECORD DELETED

## LINE, <n>

This sets the input/outout line length to <n> characters. The default is 80 characters.

## PRECISION, <n>

This sets the format used for the disolay of real values to fixed precision with <n> olaces after the decimal point.

## SIGNIFICANCE, <n>

This sets the format used for the clsplay of real values to an> significant figures. The default is 5 sigrificant figures.

## AUDIT

This command turns on the aucit faclifty lif it has been turred off previously with a NOAUDIT commanc). While it is on, al inout to the system from the INPUT flle (the terminal when running irteractivety) is echoed on the AUDIT file. Besides oroviding a recordof activity, this file may be used to update a backup copy of the catabase if a run is terminatec abncrmally le.g. a file-preserving system creshl.

NOAUDIT

This command turns off the aucit facility Normalivg it should be the first command in batch rur, since in batch mode the INFUT file may be seved and used in lieu of an aublt file.

STOP
This signats the enc of a. run. The executive will perform necessary housekeeping furctions and terminate.

### 3.3 EXECUTIVE EFROR MESSAGES

The error messages that may be generated by the executive ere summarized in this section. The prefix ***ERROR***, which is common to all messages; has been omitted.

INVALIO COMHAND
Elther the command verb was unrecognlzed, or the command syntex was incorrect.

INVALID RECORO NUMEER
The record number specified was arger than the rumber of reconds in the current set.

INVALID SET NUMEER

The set rumber specified was larger than the number of existing sets.

RECORD NOT DELETED
It was not possitle to delete the specifled record.
RECORD NOT FULLY DE-I尚DEXED
If a recoro was being deleted, it was not possible to remove same or ell of its index entries. If a record was being medified. it was not possibie to remove some or all of the index entries jor trose key data elements whose velues were changed. As a result the record mav satisfy some queries that 1 t should not.

## RECORD NOT FULLY INDEXED

If a new record was being added it was not possible to make scme or all of its index entries. If a record was belng rodified, it was not possible to rake some or all of the new index ertrles for those key data elements whose talues mere changed. As a result, the recoro will not satisfy some oueries that it shoulde

## RECORD NOT RETRIEVET

It was not possibie to retrieve a record specified in a LIST. PRINT, DUMP, MOOIFY, or DELETE command.

RECORD NOT STOREO

It was not posslbie to store a new or nocifled record.

## WORK SPACE EXCEEDED

There was not sufficient work sozce to carry out scme recuesteo operation. Eefore it wlil be possibie, the system must re recomplled with a larger work space.

### 4.1 INTRODUCTION

The ecitor is the subsysten of BDMS which is irvoked by the executlve to create a new record in core prior to its additior to a database (ADD executive ccmmand or to mocify an existing record which has been fetchec irto core from the database cmCDIfy executive command). It is also irvoked by the database definitior compller to create or modify the cefirition of a new database.

It has two operating nodes- Eppend mode, which is usec to bulld or extend a record, and edit mode, which allows selective replacement or modification of data element values and insertion or deletion of data element occurrences. Normatiy, a new record will be created in append mode, but one might shlft into edit mode to correct an error or omission before the record ls stored. Likewlse, most editirg of an existing record probably will be done in edit mode, but append mode might be used to extend the record by simply adding data element occurrences.

The editor reads free format inout comprising commands and data element names and assicnment statements. The commards tell it what oper ations to perfort on the record, thus setting the operating mode. Data element rames may occur alone or as part of an assiorment statement, depencirg on the operation specifled by the last command. In edit mode trey usualiy are qualified by one or more occurrerce numbers which identify the particular data element occirrence to be altered. Assidment statements assign values to data element occurrences.

Section 4.2 of this chapter will describe in detail the form of assignment statements to prepare the reader for the clscussion of editor commands anc operating modes in Sections 4.3-4.5. Section 4.6 15 a summary of the error messages that might be encountered while using the editor.

### 4.2 ASSIGNMENT STATEMENTS

When a value is to be assigned to a data element occurrence, it is done with an assigrment statement whose generíl form is-
<data element name> = <value>;

## data element names

A data eleffent name is elther the preferred name or one of the synonyms specilled in tre database definitlon. In edit mode, it may be qualified by ore or more occurrence numbers -- Integers separateo from the name and each other by periods. Occurrence rumbers will be discussed in Section 4.5.

## NUMERIC VALUES

The value of an INTEGER tyoe data element must te a series of digits prececec oftiorally by a + or - sign. The value may be expressed as a binary, octal, decimal, or hexadecimal rumber; the bese is specified by a sirgle letter following the number as follows-

B
Einary
or or octai
D or absent Decimal
H or 2 Hexadecinal
The set of allonable dicits is that subset of (0-9, A-F) aporocriate for the base chosen.

The value of a REAL or DOUBLE tyoe data element must te a series of decimal digits oreceded optionally by a + or - sign. it may, in addition, have a decimal coint andfor a power-of-ten multiplier expressed as an E followed by the integer exponent with an optioral sign.

Blanks anywhere within a numeric value are lgnored anc the value is terminated by a semicclon. Numeric values must not exceec the maxlmum size which car be stored as a single precision number for double orecision number for dolele type) by the computer belng used.

A numeric vector velue is represented as a series of component values constructed according to the above rules, separated by comas and terminated by a semicolon. For example,

$$
x=1.5,3.2,7.6 ;
$$

## CHARACTER STRING VALUES

The value of CHAR tyoe data element is the character string between the relational operator and a terminating semicolor. Leadirg and trailing blanks are igrired but embedded bianks are considerec to be part of the value. If leading and/or trailing blarks ere desired, they may be forced trrough use of the symol for logical negation, (a backward slash in ASCII character setl. Each leacinc or trailing negation symbol is corvertedinto a blank when the value is stored. If the character string contains semicolons, they must be doubled to avoid confusion with the terminating semicolon: any semicolon lmmediately followed by another semicolon will be sterec as a single semicolon and will not terminate the string value.

NULL VALUES

```
    A data element occurrence of any type may be assigned a null li.e. undefined) value by an assignment statement of the form
    <cata element name> = ;
This is occasionally useful if it is necessary to enter a value fon a subordinate cata element but no value is known for its parent.
```


### 4.3 EDITOR COMMANES

The editor $i s$ controlled by commands which either set its operating mode and determine how subsequent data element $n$ ames and assionment statements will be interoreted, or cause some immediete action to be taken. All commends are single letters preceded by an asterisk and followed by a blant. They will be summarized here and discussed at length in the fcllowirg two sections.

## COMMANIS THAI SET OPERATING MODE

## Append mode

*A Append data element occurrence(s). This is the cefault mode.

## Edit mode

*R Feplace date elertent valuels).
*S Substitute string(s) in data element value (s).
*D.Delete cata element occurrence(s).
*I Insert data element occurrence(s).

## COMMANCS THAT CAUSE IMMEDIATE ACTION

** Exit frcm ecitor - record will be stored or replaced by edited version.
*C Cancel edit - do not store new or altered record.
This commano simply terminates the editor and retirns control to the exective. If a new record was being input, it will not be stoned in the database. If an existing record was being editec, no changes rade. will be reflectec ir the databese. It is prlmarliy useful in an interactive session to terminate a hopelessly confounded edit or one which was mistakenty begun.
*L List the present state of the record.
This commanc allcus a record to be visually checked anytime while it ls being created or edited. It does not affect the coerating mode or atter the edit command in effect when it is civen.
*e End of record.
This marks the erc of a record when multlole records are being

ADDed. It is usualiy not needed (c.f. section 4.4, subsectlon on automatic record generation). When MODIFYirg a record, it has the same effect as **.

### 4.4 APPEND MODE

When first entered, the editor ls in appenomode. It may be
returned to append mode ef any time with a *A commard.

## ORDER OF ASSI ENMENT STATEMENTS

In apoend mode the order in which assignment staterents and pure node data element names are entered determines the order of the deta element occurrences crested and their linkage to parent occurrences. The occurrences of a data element Inked to each occurrence of its oarent cata element for tre recorc level data elementl form an ordered IIst. In general, each essignmert statement or pure node name creates a new occurrence of a data element which wlll be addec to the end of the list of occurrences that is linked to the last=created occurrence of its parent data element.

For examole, if data were belng added to the databese descrited in Chapter 1.
$A U T H O R S * ;$
$A_{\bullet}=J O N E S ;$
$A_{\bullet}=S M I T H:$
$I_{\bullet}=L B L ;$
AUTHORS*
$A_{\bullet}=E A K E R ;$
$I \bullet=U C B ;$
would create two AUTHORS groups, the first with two authors at LEL and the second whose sirgle zuthor is affilateo with UCE.

The periods following the cata element names mean rext*. A data etement name without a period means *irst* or occurrence number 1. The next two subsections will elucidate the signfflcerce of these conventions.

Other than havinc cata element occurrences in the cesirecorder and following the correct parent data element occurrence, the order of assignmert statements ir appenc mode ls immaterial An assicrment statement for the reference, $R$ could have been insertec anywbere in the preceding example nithout effecting the creation of the AUTHCRS groups.

## AUIOMATIC PARENT NODE GENERATION

In apoend mode, $N O D E$ type parent data element names cen usualiy be omitted from the lrput stream. The occurrences of suct data etements necessary to lirk togetter subordinate data element occurrences will be generated cutomaticaliy if

1) One of the suborcinate deta elements is encountered ir the input stream and no parent occurrence yet exists, or .
2) occurrence number 1 of one of the subordinate cata elements is enccuntered, l.e. the data element name apoears without a period following it, and the lest-created parent occurrence already has at least one occurrence of that cata element linked to it.

Thus, the prececirg example could heve been entered ir the simpler form
$A=J O N E S$;
A. $=$ SMITH;
$I=L E L$;
$A=E A K E R$;
I=UCE;
or eoulvalentiy,
$I=L E L ;$
$A=$ JONES;
A. $=$ SMITH;

I=UCE;
$A=E A K E R$;
Each occurrence of $A$ in the first case or I in the secord causes the creation of an AUTHORS node.

Automatic parent node generation will also work if the parent data element is not NCDE type. In this case, an automatically generated parent occurrence will have a null value.

## AUTOMATIC RECORO GENERATION

When ADDing records, if a record-level data element name appears more than once without a fllowirg perlod, the second occurrerce will trigger the storace of the previous record and begin a net one just as if a *E command hac been encountered. The second occurrence of the record-level data element wlll become a oart of the new reconde

The record itsell may be conslderec to be the pure rode perent of all record-level cata elements, so automatic recoro cenenation is comoletely analogous to the automatic parent node generation described in the oreceding subsectlon.

If the eoltor has been invoked by a monify executive command, automatic record generation is meaningless, so that in acpend mode, a recurrence of a record-level cata element name without a peniod will be flagged as ar inout error.

## MULTIPLE ASSIGNMENT STATENENTS

Several consecutive cccurrences of a data element may be created without the necessity of repating the data element rame. ith is done with a multiple assignment statement, whose most gereral form is
<CE name> = <value $1>$;<value 2$\rangle ;$....
Using a mutiple assignment statement, the oreceding examole could have been inout in the still simpler form

$$
\begin{aligned}
& A=\text { JONES; } \leq M J T H ; I=L B L ; \\
& A=B A K E R ; \quad I=U C B ;
\end{aligned}
$$

or equivalentiv,
$I=L E L ; A=J O I E S ; S M I T H ;$
$I=U C B ; A=E A K E R ;$

The differ ence between the two assignment statements

$$
x=1,2,3 ;
$$

and

$$
x=1 ; 2 ; 3 ;
$$

should be clearly urderstood. The first creates a sirgle occurrence of $X$ whose value is a three-comconent vector lassuming $X$ is defined to be a numeric type data elementl. The second creates three occurrences of $X$ whose values are 1,2 , and 3 , respectively.

In order to inplement multiole assignment statements, tre inout processor, wilt Ettemct to interoret anything which is not a recognizable deta element name or esslgnment statement as another value in the creceding assignment statement. If a cata element name is misspelled, the erroneous assignment statement cerrot be distinguished from anotrer value for a oreceding CHAR type deta element. It will only be recognized as an error if the precedirg deta element is not CHAR type.

### 4.5 EDIT MODE

Ary of the commands *R, *S, *D, or *I olaces the ecitor in ecit mode.

## NAMING DATA ELEMENT OCCURRENCES

In edit mooe, tte data element occurrence to be altered is identified by specifying its cosition within the record structure. To comoletely identify a particular data element occurrerce, one must specify its name ard occurrence number, the number of the carent occurrence to which it is lirked, the number of the granoparent occurrence to which its carent is linked, etc. up to the record-level ancestor. The serles of occurrence rumbers is appendec to the deta element name in this order, i.e., in order of lncreasing remoteress of the ancestor, as a series of intecers seoarated from the rame and each other by periods.

Referring to the example record structure of Chapter 1, A. 4.2 would identify the fourth author A. 4 of the second atthors group, AUTHCRS.2. The (first and oniv) $x$ value in the seconc DATUM in the first taELE would be identified as X.1.2.1, etc.

Under certaln circumstances it is not necessary to specify all these occurrence numbers since tre editor will supply default values. This will be discussec below in the subsection on path memory.

## REPLACE COMMAND (*R)

Tris is followed by assignmert statements which essien new values to existing dati element occurrences. For example,
*R A.2.1=J.COE;
would cause the value of A.2.1 to be replaced by J.OCE.
It is not meaningful tc repace a pure NODE type data element since it carries no value.

## SUESTITUTE COMMAND (*S)

This allows the replacement of a selected substring mithin a CHAR type data element value. It works like the 7 R command excent that the data element values within assignment statements whict follow a *s command have the form
<delimiter><cid subsiring><dellmiter>enew substring>edelimiter>e
The cdellmiter> may be any character that doas not occur in elther the <old substring> or the <new substring>. All blankswithin the substrings are significant.

The snew substring> will rectace the first substring in the data element value trat matches the cold substring>. If the cold substring> $l$ s null, the knew substring> will be irserted et the beginning of data element value. If the enew substring> is null, the first substrinc ir the data element value that matches the cold substring> will be celetec.

For example, if A.2.1 has the value J.DEO, the command *S A.z.1=,EC/OE/;
would charge it to J. TOE. The command
*S A.2.1 $=1 / 8 /$;
would change thet value to BJ. DOE, while the command
*S A. $2 \cdot 1=/ \mathrm{J} / /$;
would leave as a final velue $\mathrm{B} . \mathrm{DOE}$.
The modified data element value is echoed following execulior of a *S commend so thet the correctness of the substitution maybe ascertained.

DELETE COMMAND (*D)

This is followed by the rames of the dita element occurrences to be deleted, each terminated by a semicolon. All cccurrences of subordinate cata elements linked to a deleted data elemert occurrence are also deletec.

For example,

$$
\text { *D DATUN. } \mathrm{E} \cdot 1 \text {; }
$$

would defete DATUM.e in TAELE. 1 and any occurrences of $x, Y$, $E Y$, and co linked to it.

After a cata element occurrence has been deleted, any occurrences of that data element linked to the same parent cccurrence and following the deleted occurrence will be ldentified by occurrence numbers one smaller tran originally.

## INSERT COMMAND (*I)

This is follover by assignment statements for occurrence names for pure node tyoe date elements). Each of them will cause an occurrence of a cata efement to be inserted before the ramec occurrence.

For example, if the author name which should have been A. 2 in AUTHORS. 1 had been mistakenty omitted from a record, it could be inserted with the command

* 1 A. $\overline{\text {. }} 1=\mathrm{J}$. COE;


#### Abstract

After a date elemert occurrence has been inserted, ary occurrences of that data elenent linked to the same parent occurrence and following the inserted occurrence will be identified by occurrerce numbers one larger thar origirally. Hence the inserted occurrence gecomes the $n$ amed occcrence.


In order to add a data element occurrence after all existing occurrences lirked to a siven parent occurrence, one either specifles an occurrence number greater than the number of occurrences already existing, or specifles occurrence number 0 , or simply omits the iest qualifier, but not its precsgirg period. Hence, the command in the preceding example would rave the desired effect even 1 f duthofs. 1 had only one subordirate $A$ occurrence orior to tre insertion.

A multiole assignment statemert following an insert command causes several new occurrences to be inserted before the soecified occurrence of the data $\in$ ement. They will occur in the modified record in the order of their appearance in the assignment statement. After insertion, any occurrences of that cata element linkec to the seme parent occurrence ard following the inserted occurrences wlll heve their occurrence numbers incremented by the number of inserted occurrences.

For example,
*I A.4.2 = J.SMITH; B.JOMES;
would leave AUTHORS. 2 with $A .4=J . S M I T H$; and $A .5=B$.JCNES: If an A. 4 existed prior to insertion, it would now be A.6. etc.

## ORDER OF EDITS

One point rade in the preceaing discussion warrants furtrer emphasis - data element occurrence numbers are retative list oositicns and can change during eciting. After data element occurrences are inserted or deletec, ary occurrences following the insertion or deletion point (linked to the same parent occurrence) must be identified by occorrence numbers larger or smaller thar origirally.

In order to avoic confusion, especially when working from a orinted listing or dcing bitch ucdates, replacements and substitutions should be done first, ald insertions and deletions should be cone from the bottom uo, l.e. these data element occurrences witr the tighest occurrence numbers shculd be edited first. Then successively modifled occurrences will retain tre occurrence rumbers of the original record.

If one does become confused while working irteractively, the current values of the occurrence rumbers may be ascertaired by listing the record with a* command.

## PATH MEMORY

In edit mode, the ecitor will attempt to supply a default parent hierarchy for a data element beling modifled if the user has not giver complete dosition irfcrmetion.

This is implemented with a path memory, which works as follows. The editor renembers the path taken through the record structure to arrive at the position of a modification. Then, if the position of the next modification is not completely specifled, an attempt will be made to link whatever partial path is specifled to the ofth remermered from the last modilicetion. If thls linkage can be echievec, the resulting path will be used as the position of the new modification and the path fermory updated.

The path memory is cleared when the editor enters appenc mode.
The best way to understand the effect of this generzi mechanism is through the consideration of some specilic examples.
1.: Suppose one wishes to replace the values of several A"s within AUTHOFS.2. This can be accomolished with the cormard

$$
\begin{gathered}
\text { *R A. } 1 \cdot 2=\text { E.JONES; } \\
A \cdot 2=E \text { SMITH; } \\
A \cdot 4=C \text { EAKER; }
\end{gathered}
$$

The occurrence of the parent data element AUTHORS is not specifled for $A .2$ and $A .4$ and 50 it is assumed to be AUTHORS.2, as specified for A. 1 .
2. Suppose that the entire second AUTHORS group was inadvertentiy skipoed when a record was initialiy input. If could be inserted with the command

```
*I AUTHORS.2;
    A=E.JCNES; E.SMITH; J.DOE; C.BAKER;
    I=LBL;
```

Since no parent occurrence is specifled for A.1. A. 2, A.3, A.4, and I. 1, it is ploked up from the path memory as

AUTHORS.2.
3. Suppose that a new third datum is to be adceo to TAELE. . $^{\text {. }}$ It can be inserted with the command
*I DATUM.3.2;
$X=1.537$;
$Y=3.20 € ; \quad$ $Y=0.001$;
CD=NEW DATA POINTS:
Since no pirent occurrence is specified for $X .1, Y, 1$, and CD.1, it is cicked up from the path memory as DATUM. 3 Ir TABLE.2. For DY.1, it is assumed to be Y. 1 in CAJUM. 3 in TABLE.2.
4. Suppose that after the edit in $3 .$, one deslres to replece the value of $X .1$ in $D A T U M .2$ in TABLE. 2 . The commanc
*R X $\quad 1.2=2.372 ;$
will accomplish this. It is not necessary to give a complete oosition soecification for $X$, namely $X .1 .2 .2$, because the default occurrence of TAELE resulting from the crevious edit is TABLE. 2 .

### 4.6 EDITOR ERROF MESSAGES

In this section, the ecitor error messages will be summarized and their meanings elaborated. It must be clearly recognized, however, that the editor is designed to make every effort to find a legal interpretatior for all input it processes. The lmplications of tris were discussed in section 4.4. (subsection on multicle assignment statements).

In gereral, when an error is detected, that oert of tre incut stream which caused the error will be skipoed and processing will continue with the following incut. In append mode, tris means trat some data element occurrence will not be created. In edit mode, scme edit operatior will not be performed. If one is working interactively, the error usually can be corrected immediately. In batch mode, one error will sometimes cause several cthers, e.g. a parent data element occurrence is not created due to a misspelling, causing severa: scbordinate cata element occurrences to be linked incorrectiy to a previousiy created parent occurrence.

All error messages issued by the editor are prececed by ***ERROR***. This prefix has teen omitted in the follcwing list of messages.

## data element value in delete command

It is not mearirgful to specify a new value for a data element occurrence being celetec. Onty data element names terminated by semicolons may follow a delete command.

## Illegal recurrence cf record-level data elenent

A record-level data efement name without a perico occurred in append mode while MCOIfying a record which already ras en occurrence of that data element. While ADBing records; this is nct an error but will cause the cutomatic generation of new record.

## INCOMPLETE ASSIENMENT STATEMENT

In batch mode, an end-of-ilile condition has betr sensed while processing a cata elefent name or value. Probably tre teritirating semlcolon was omitted.

## INCORRECT QUALIFIER

The editor could nct determine the oath to a data element occurrence basec on the occurrence numbers given and the current state of the path memory. Thls message is also given lif toc mary occurrence numbers are specilled for the depth of a data elefert within the hierarchical record structure.

INVALID LENGTH RETURNED EY IPROC
The user-sucplied Iffoc routine has returned a regative data element length.

## INVALID NUMERIC DATA ELEMENT VALUE

The value specifled for a numerlc data element does not have $\bar{f}$ form valio for its type.
missing data element value
A non-NODE tyoe dafa elemert name is followed imrediately by a
semicolon. this ls only allowed in a defete comand.

MULTIPLE ASSIGNMENT STATEMENT IN REPLACE OR SUESTITUTE COMMAND
Only one data element occurrence at a time may be reolaced. Multiple assignment statements are not meaningful following a reolece or substitute command.

NO MATCHING SUBSTRING
The old substring specifled in a substitute commanc does not occur in the strirg being scanred.

NON-CHAR DATA ELEMEAT IN SUBSTITUTE COMMAND
A substitute comrand may onfy operate on CHAR tyof deta elefents. PARENT DATA ELENENT MISSING

An edit operation cannot be performed because an indicated oarent data element occurrence coes not exist.

PURE NODE DATA ELEMENT IN REPLACE OR SUESTITUTE CCMMAND
Since a NCDE type data element carries no value, it is mearirgless to replace lt. Only assignmert statements may follow a reflace or substitute command.

QUALIfIE data element name in apfend mode
Deta element nemes may not be followed by occurrerce numbers in apperd mode.

## SPECIFIED OCCURFENCE LOES NOT EXIST

The data element occurrence specified in an edit colmand does rot exist. An odcurrence number probably was inout incorrectly.

## SUBSIITUTE COMMAND SYNTAX

The syntax of a substitute cownand is incorrect, e.g. the delimiter character does not occur exactly three times or there are non-blank characters between the trird dellimiter and the terminatirg seficolcr.
undefined data element name
Tre input stream contains a data efement name which coulc not be

```
recognized and covid not be Interpreted as another value ir a
preceding assigrment statement.
```


## HORK SPACE EXCEEDED

A deta elenent name or value is longer than the evallable work space or there is no more, spece avaliatie for expansicr of the peth - memory. In general; the system must berecomplied with a longer work space to guarentee that the errer will not recur when accessing this database. Sometimes, storing the record and then retrieving it agein will free enough work space to alleviate the problem anc allow further editing. This is likely lif extenslve edits have resulted ln a large amount of deac space in the recorc.

Certain error coiditions that generate messages car arlse while tisting a recond They are sumarized in this secticr.

A11 error messages issued by the list routine are preceded by .***ERROR***. This prefix has been omitted in the folicwing list of messages.

INVALID LENGTH GETURNED BY OPROC - DATA ELEMENT SKIPPED
The user-sucpiled opROC routine has returned a negative data element length.

INYALID MODE PARAMETET - LIST ARCRTED

This wili cniy occur when the user calis LIST directig with an invalid mode parameter.

## RECORD BUFFER EMPTY -- LIST ABORTED

There is no record in core to be isted, not even an empty one. This will oniy occur if a user calis LTST directiv without croperiy initiallzing the recoro butfer.

WORK SPACE EXCEEDED - DATA ELENENT SKIPPED

WORK SPACE EXCEEDED -- LIST ABORTED
The record is sc arge that there is insufficient work soce for LIST to operate. The system should probably be recomplled witr a larger work space for use with thls database.

### 5.1 INTRODUCTION

The BOMS auery language permits a user to search a datatase for those records satisfying an arbitrarily complex concition on key (indexed) cata element values. The cordition is constructed as a Boolean combination of key value specificaticns, includino Inequalities and ranges. Furthermore, it is possible to search for records having an occurrence of a specifled data element regardiess of value, or for those having an cccurrence of the cata element with a null value. Truncated value sceciflcation for character string keys may be used to search for those records having an occurrence of the data element beginning in a particular way.

The retrievai facility is linvoked by the FIND execulive command. This must be followed by a condition that is terminated by a*, i.e.

FIND <conditior* **
The set of recorcs that satisfy the condition is essigned a set identification number, and the number of records ir the set is reported in the form.
(n) RECORD (S) IN SET <S>

If no records setisty the condition, the response ls
0 RECORD(S) IN SET
The number <s> assisned to a non-empty set may be usec in subsequerit aueries in combination with further concitions.

The rext two sections describe the format of conditions.

### 5.2 SIMPLE CONOITIONS

A simple concition has the form -

where curly brackets [] surround a set of options, one of which must be chosen, and square brackets $[1$ surround a completely ootioral element. The relational operators $\ll>=, \ll$ stand for liess than or equal*. "greater than or equal*, and not equal*, respectively. Thus, an exact value, an.incleslve or exclusive upper or loner bound, or a range of values may be scecifled.

Examples of valic simple concitions are -
$A=$ JONES;
A <> SMITH:
$x<=7 ;$
$x=5 ;$ to $8 ;$
$X>5 ;$ TO<8;
The later two examoles differ in that the secondexcludes both endpoints of the range.

The condition
<key data element name> <> <VALUE>;
will be satisfied by all records which
a) have at least one occurrence of the named data elentent, and
b) have no occurrence of that data element whose value matches that glven.

If no value appears after the relational operators = or <>, a search is made for those records which respectively do or do not have an cccurrence of that data element with a null value. Note trat this corresponds to the syntax used to enter a null value usirg the editor (cf. Chapter 4). For example,
$A=;$
would result in a search for all records having a null velue for $A$ and
A <>;
defines the complement set.
The last for allowed for a simple condition, e.g.
<key data el ement names:
defines the set of aif recor os maving any occurrence of the specifled data element, regarcless of value.

The format of a value depends on the data element type and ls the same as that acceptea by the editor (cf. Chapter 4). Key values in an index are fixed length, derived from the corresponding data element values by truncatior or padding; data element values in a query will be truncated or padded in the same way to ensure valld comparison with index values.

### 5.3 COMPLEX CONOIIIINS

The most general condition that may apoear in a ouery is constructed out of simple condtions and previousiy defined $s \in t s$ according to the lollowirg recursive definition -
<CONOITION> $=\{N O T]$ \{ SIMPLE CONDIIION>\} \{ [AND\} <CONDITION>]
C<SET NLMBER, \} [ COR 3 ]

That ls, simple conditiors and oreviously-defined sets, identifled by number, may be combinec using the Boolean operators NOT, AND, and OR. NOT has the hlghest precedence and OR the lowest; this ordering may be overridden trough use of parentheses.

Note that simple concitions and set numbers olay eoulvalent roles In a complex query. This is because each simple condition may be viewed as defiring ar (irtermediate) set. These sets, elong with any existing sets appearing in the query, are then comblned by unlon (OR), intersectior (AND), and complement (NOT).

It should be noted that the tmo conditions <key data element name> <> <VALUE>;
and

```
NOT <key deta \epsilonlement name> = <VALUE>;
```

are not equivalent. The meanirg of the first was elucidated in the oreceding section. Tre seconc differs in that records mevirg no occurrence of the namec data element will also satisfy lt. Tris happens becuase NOT complements the set defined by aflowing condition and the universal set in terms of which cormolementation is defined is the entire catabase. Thus, the effect of a Not operator cannot simply be absorbed into the relational operator and
NOT $=$ <>
NOT < $>=$
NOT $>$ are not equivalent to $<=$
NOT <= >
NOT $>=$ e
NOT <> =

### 5.4 RECORD ID

As mentloned ir Chapter 1, every record in a database is assigrec a uniave and permanent integer record ID when it is first created. The pseudo cata element neme REC-ID way be used in a simple condition to retrleve records oy their ID*. Any form of the simole condition including ranges and

REC-ID:
is allowed. The set of records satisfying the latter condition* is the entire database.

Simple conditions en: REC-ID maz be used as components of a complex condition.

### 5.5 TRUNCATED SEARCH

Truncation provides a way of searching for records that contain occurrences of some data element whose value begins ir a specilied way, regaraless of the rest of the value. for examole, cremight like to ind all papers authored by Smith, regardless of how fre first neme had been entered ir. the catabase. If authors names were stored iest name first, one could just do a truncated search for Smith. Truncatior is indicatec in the BOMS query language by a single/ (slash) following the cartial value. Thus, the preceding example would be expressec as

FIND $A=S M I T H /$; **
Any record with a value for A whose flirst five characters are SMITH would satisfy the concition.

If one desires to search for a value which happens to end witr a/, the slast may be doubled to prevent its belng interoreted as a truncation dellmiter. For example,

FIND CT=ABC//; **
would find all records having a value $A B C /$ for $C T$. If cre desires to truncate immediately following a/, then three slashes are reaulred. For example,

FIND CT=ABC///; **
would find all records containing values of $C$ which $t \in g i n$ with the characters ABC/.

Truncated values may be used in conjunction with ary relatioral operator. Elther or both of the endpoints of a range search may be expressed as truncated values in the same wav. Since equality to a truncated value actally defines a range of values that will setisfy the concition, the meanincs of the remaining relational ocerators will be affected in a corresponding way, elucidated by the following diagram -


The range of record values setisfying an equality condition is represented by the middle segment of the top iline whlle the ranges satisfying other corditions involving the same value but different relational operators are represented by the other line segments.
simple conditions involving truncated values may be useo as components of a comflex concition.

### 5.6 RETRIEVAL EFROR MESSAGES

The error messeges that can be generated whlle processing a auery are listed below. The ertire query is scanned for errors and only if none are detected is the catabase actualiy searched.

All error messeges from the query processor are preceded by ***ERROR***. This prefix has teen omitted in the following list of messages.

## INCOMPLETE QUERY

In batch made, an end of file condition has been sensed whlle processing a query. Probably tre query terminator (**) was omitted.

INVALID LENGTH RETURNED EY QPROC
The user-sumplied QFROC routine has returned a regative data element length.

INVALID NUMERIC DATA ELEMENT VALUE
The value specified for a numeric data element does not have a ferm valld for its tyoe.

## INVALID RANGE SFECIIICATION

A meaningless set of relational operators has been used in a rarge condition, e.g. $X<7 ;$ TO $>$ 9;.

INVALID SET NUMEER
A set number has been usec that is less than or ecual to o, or larger than the last-created set number.

NON-KEY DATA ELEMENT IN QUERY
A data efement used in the auery is not defined to be a KEY. QUERY SYNTAX

The query has not been proderiy constructed cut of simele conditions, and Boclean operators, e.g. two or more simple conditions are not joined by a Boolean operator.

TOO MANY LEFT PARENTHESES
TOO MANY RIGHT FAREITHESES
The parentheses appearing ir the query are not properiy paired. UNDEFINED DATA ELEMENT NAME

A data element name appearing in the query cannot be recognized. It has probably been misspelled.

## VALUE SPECIFIED FCR PURE NODE DATA ELEMENT

The query includes a condition on the value of a NODE data element. This is meanirgless since a pure NODE carries no value. Perhaps the data element rame is misscelled.

## WORK SPACE EXCEEDED

The system has run out of work space while processing the ouery. It may be necessary to recomplle the system with a larcer work space for use with this detabase.
6.1 INTRODUCTION

Utility programs are provided with goms to perform the mairterance functions of initial database loaing, dumping an entire database in external format for transmittelt compressing dead space from the deta file, and rebalancing the index trees. Use of these programs is explained in the following sectiors.

### 6.2 LOAD

The $L O A D$ utitity is used for initlal batch loading of a database. At present, it cannot be used to edd records to an existirg database that must be cone with the executive ADD command LOADing ls a much more efficiert operatior than ADDing cecause the index entries are saved until all records have been stored in the cata flife, and then the index trees are built from the bottom up in a single pass. Thus, the index trees are perfectly balanced, leading to maximum efflclericy in query processing.

Inout to the LOAC utility consists of one or more commands flioned by the data to be lcaded in exterrai editor format, with records separated by $* E$ and tre lest record terminated by **. This is the format produced by the DUMP utility or the executive DUMP conmard. The commands that may precede the input data are -

LINE, <n>
This sets the input line length to un characters. if no LINE command is present, the line length defaults to 80 characters.

FREE; <n
This sets the amount of free space to be eft on each index page to allow future catabase expansior without the possibie consequerce of unbalancing the index trees. If kns <s o, <n> percent free soace will be left on each page. If <n> = 0, the pages will be comoletely fliled to achieve minimum index slze for a static database.

LOAD
This terminates the command stream and initiates the data stream. A LOAD command mist precede tre data even lif no other commands are oresent in the incul stream.

These commancs are terminated by a blank and hence must not contair embecded blanks.
6.3 DUMP

The DUMP utility is used to dump an entire database in format sultable for subsequent reloadinc. This ls prlmarily useful if it is necessary to transmit the database to another site. To achleve maximum efficiency, the records are outout in the (rardom) order in whlch they occur in tre data file, rather than in record Io or key seavence. If it is desired to dump only selected oarts of a detabase, this may be done with the executive FIND and DUMP commans.

The DUMP utility reads from the input file commands scecifyirg the format of the dumb file. They are -

## LINE,《"n>

This sets the ottout ine length to en> characters. If no LINE command is present, the line length defaults to eo charectens.

## EXPAND

This selects an expanded dump format in which each data element occurrence begins a new line. This may be useful if a texteditor is to be used tc modify the dump flle prior to reloading the database. If no expand commard. is oresent, a more compact and efficientiy readable format is usec in which the outout lines are completelv filled.

These commands are terminated by blank and nence must not contair embedded blanks.

## CLEAN

### 6.4 CLEAN

The CLEAN utility is usec to compress out of the data flle the dead space resulting from update activity. The frequency with which this operetion shoulo be cerried out depends on the nature and frequency of updete activity for the datzbase.

### 6.5 BALANCE

The BALANCE utility is usec to rebalance the incex trees. This operation may be necessary after extensive update activity in order to ootimize auery processing.

The input file optionally may contain the folloting commard. FREE, <n)

This sets the amourt of free space to be feft on each index page to al low future cetabase expensior without the possible conseduerce of unbalancing the irdex trees. If <n> < 0 , <n> percent free space will be left on each page. If $\langle n\rangle=0$ the pages wll be completely fllled to achieve minimum incex size for a statlc database.

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