EFFICIENT OPERATOR-COARSENING MULTIGRID SCHEMES FOR LOCAL DISCONTINUOUS GALERKIN METHODS

DANIEL FORTUNATO*, CHRIS H. RY CroFT*, AND ROBERT SAYE†

Abstract. An efficient $hp$-multigrid scheme is presented for local discontinuous Galerkin (LDG) discretizations of elliptic problems, formulated around the idea of separately coarsening the underlying discrete gradient and divergence operators. We show that traditional multigrid coarsening of the primal formulation leads to poor and suboptimal multigrid performance, whereas coarsening of the flux formulation leads to optimal convergence and is equivalent to a purely geometric multigrid method. The resulting operator-coarsening schemes do not require the entire mesh hierarchy to be explicitly built, thereby obviating the need to compute quadrature rules, lifting operators, and other mesh-related quantities on coarse meshes. We show that good multigrid convergence rates are achieved in a variety of numerical tests on 2D and 3D uniform and adaptive Cartesian grids, as well as for curved domains using implicitly defined meshes and for multi-phase elliptic interface problems with complex geometry. Extension to non-LDG discretizations is briefly discussed.

Key words. discontinuous Galerkin methods, multigrid methods, elliptic interface problems, implicitly defined meshes

AMS subject classifications. 65N55, 65N30, 65F08

1. Introduction. Discontinuous Galerkin (DG) methods have gained broad popularity in recent years. They are well-suited to $hp$–adaptivity, provide high-order accuracy, and can be applied to a wide range of problems on complex geometries with unstructured meshes. Although DG methods were first applied to the discretization of hyperbolic conservation laws, they have recently been extended to handle elliptic problems and diffusive operators in a unified framework [4]. Such methods include the symmetric interior penalty (SIP) method [16, 3], the Bassi–Rebay (BR1, BR2) methods [6, 7], the local discontinuous Galerkin (LDG) method [15], the compact discontinuous Galerkin (CDG) method [25], and the hybridizable discontinuous Galerkin (HDG) method [13]. In particular, the development of efficient solvers for DG discretizations of elliptic problems is an active area of research.

Among the panoply of DG methods for elliptic problems, the LDG method is a popular choice: it is accurate, stable, simple to implement, and extendable to higher-order derivatives [35]. Additionally, on Cartesian grids it has been shown to be superconvergent [14]. The LDG method results in symmetric positive (semi)definite discretizations which are well-suited to solution by efficient iterative methods. In particular, the multigrid method has emerged as a natural candidate due to its success in the continuous finite element and finite difference communities, both as a standalone solver and as a preconditioner for the conjugate gradient (PCG) method. However, direct application of standard multigrid techniques to DG discretizations of elliptic problems can result in suboptimal performance [1, 19], and much work has gone into developing specialized smoothers and coarse-correction methods to remedy this issue,
even on Cartesian grids [21, 17].

When a mesh hierarchy is available, geometric $h$–multigrid is a natural choice of solver. Error estimates have been derived for a multilevel interior penalty (IP) method on unstructured meshes, yielding convergence factors in the range $\rho \approx 0.3$–0.5 for Poisson’s equation [19, 9]; the method has also been applied to adaptively refined Cartesian grids with similar results [22]. Subsequent work describes how the multilevel IP method [19] can be used as a preconditioner for the LDG method—both for the Schur complement system (using conjugate gradient (CG)) and for the saddle-point system (using GMRES)—resulting in a bounded condition number with respect to mesh size $h$ [21]. More recent work for LDG and IP uses a multigrid W-cycle on nested [1] and agglomerated [2] unstructured meshes; however, these results indicate poor convergence factors of $\rho \approx 0.8$–0.9 even with many smoothing steps.

A popular choice for high-order DG methods is $p$– or $hp$–multigrid [20, 18, 23, 5], where $p$ refers to coarsening the polynomial degree $p$ in the multilevel hierarchy, and $hp$ refers to some combined strategy of coarsening the mesh size $h$ as well as the polynomial degree $p$ of the underlying discretization. Using factor-of-two coarsening in $p$ with an element Jacobi smoother, convergence factors of $\rho \approx 0.5$ were achieved for Laplace’s equation with $p \leq 4$ [20]. Fidkowski et al. [18] used a line smoother with sequential coarsening in $p$ that gave similar results for convection–diffusion problems, though the performance degraded as $h \to 0$. A method employing an overlapping Schwarz smoother with factor-of-two coarsening was used with PCG on LDG discretizations up to $p = 32$ [31]; this method exhibited good convergence factors of $\rho \leq 0.1$ on high aspect-ratio Cartesian grids, at the cost of an expensive smoother.

Algebraic multigrid methods have also been applied to DG discretizations of elliptic problems. A hierarchy of operators can be defined by agglomerating neighboring unknowns based on smoothed aggregation (SA), resulting in average convergence factors of $\rho \approx 0.4$ and $\rho \approx 0.2$ for the bilinear BR2 and SIP methods, respectively [26]. An SA method employing energy minimization was used with PCG to achieve $h$-independent convergence factors of $\rho \approx 0.2$ for LDG discretizations, but performance degraded with increasing $p$ [24]. A method based on unsmoothed aggregation was developed for the IP method using a coarse space consisting of continuous linear basis functions [8]; this method proved robust for multi-phase problems with large jumps in ellipticity coefficient, but efficiency weakly degraded with mesh size. A related approach based on smoothed aggregation and low-order coarse correction yielded similar results [30].

Independent of the type of multigrid method, a particular fact to note—and something we believe underpins the difficulties in applying multigrid to DG—is that coarsening a fine-grid operator is not always the same as constructing that operator directly from the coarse grid. Indeed, it was noted by Antonietti et al. [1] that, for the LDG method, “convergence cannot be independent of the number of levels if inherited bilinear forms are considered (i.e., the coarse solvers are the restriction of the stiffness matrix constructed on the finest grid).” In this paper, we analyze why this is the case and show that this facet can be remedied if one coarsens with care. In particular, we present a simple modification to traditional multigrid operator coarsening that yields optimal multigrid convergence for LDG (and can be extended to other DG discretizations of elliptic problems). We show that traditional coarsening of the fine-mesh elliptic operator results in poor performance, whereas coarsening of the saddle-point flux formation restores optimal multigrid efficiency. Our approach is equivalent to pure geometric multigrid but avoids the need to explicitly build the coarse mesh and its associated components, such as quadrature rules, Jacobian mappings,
lifting operators, and face-to-element enumerations—as discussed, this holds benefit for a variety of intricate DG implementations where building the coarse mesh can be problematic.

The paper is structured as follows. In section 2, we formulate a general DG discretization of Poisson’s equation and derive the LDG method through the appropriate choice of numerical flux. In section 3, we describe the construction of geometric $hp$–multigrid methods in the corresponding DG setting. In particular, we show that traditional operator coarsening can fail to create the true coarse operator for LDG discretizations, and present a modified coarsening strategy that faithfully reproduces the coarse operator. In section 4, we present numerical results for the standard and modified multigrid methods on uniform and adaptively refined Cartesian grids in 2D and 3D. We conclude with some examples of multi-phase elliptic interface problems on implicitly defined meshes, which demonstrate good multigrid performance even on meshes with long and thin filaments as well as tiny and dispersed phase components.

2. Discontinuous Galerkin formulation.

2.1. Model problem. The model elliptic problem considered in this work is the Poisson problem

\begin{equation}
-\nabla^2 u = f \quad \text{in } \Omega,
\end{equation}

\begin{equation}
\begin{align*}
u &= g \quad \text{on } \Gamma_D, \\
\nabla u \cdot n &= h \quad \text{on } \Gamma_N,
\end{align*}
\end{equation}

where $\Omega$ is a domain in $\mathbb{R}^d$, $\Gamma_D$ and $\Gamma_N$ denote the components of $\partial \Omega$ on which Dirichlet and Neumann boundary conditions are imposed, $n$ is the outward unit normal to the boundary, and $f$, $g$, and $h$ are given functions defined on $\Omega$ and its boundary.

2.2. DG for elliptic problems. In order to apply a DG method to (2.1), we rewrite it as a first-order system by introducing the auxiliary variable $q = \nabla u$ and writing the Laplacian as the divergence of $q$ [4]:

\begin{equation}
\begin{align*}
q &= \nabla u \quad \text{in } \Omega, \\
-\nabla \cdot q &= f \quad \text{in } \Omega, \\
u &= g \quad \text{on } \Gamma_D, \\
q \cdot n &= h \quad \text{on } \Gamma_N.
\end{align*}
\end{equation}

In this work, we mainly consider discretizations of (2.2) wherein the corresponding meshes arise from Cartesian grids, quad/octrees, or implicitly defined meshes of more complex curved domains (see sections 4.1, 4.2, and 4.3, respectively). As such, it is natural to adopt a tensor product piecewise polynomial space. Let $\mathcal{E} = \bigcup E$ denote the set of elements of a mesh of $\Omega$, let $p \geq 1$ be an integer, and define $Q_p(E)$ to be the space of tensor product polynomials of degree $p$ on the element $E$. For example, $Q_3$ is the space of bicubic (in 2D) or tricubic (in 3D) polynomials having 16 or 64 degrees of freedom, respectively. We define the corresponding spaces of discontinuous piecewise polynomials and vector fields on the mesh as

\begin{equation}
V_h(\mathcal{E}) = \{ v : \Omega \to \mathbb{R} \mid v|_E \in Q_p(E) \text{ for every } E \in \mathcal{E} \},
\end{equation}

\begin{equation}
V_d(\mathcal{E}) = \{ \omega : \Omega \to \mathbb{R}^d \mid \omega|_E \in [Q_p(E)]^d \text{ for every } E \in \mathcal{E} \},
\end{equation}

respectively. We denote by $(\cdot, \cdot)$ the natural $L^2$ inner product on $V_h$ and by $\| \cdot \|$ the corresponding norm, $\|u\|^2 = (u, u)$, with analogous definitions for $V_d$. 

In a DG method, both \( q \) and its divergence are defined weakly via numerical fluxes defined on each mesh face. The weak form of (2.2) consists of finding \((q_h, u_h) \in V^d_h \times V_h\) such that

\[
\int_E q_h \cdot \omega = - \int_E u_h \nabla \cdot \omega + \int_{\partial E} \tilde{u}_h \cdot \omega \cdot n, \tag{2.5}
\]

\[
\int_E q_h \cdot \nabla v - \int_{\partial E} \tilde{q}_h v \cdot n = \int_E f v, \tag{2.6}
\]

for all test functions \((\omega, v) \in [Q_p(E)]^d \times Q_p(E)\) and for all \( E \in \mathcal{E} \). The numerical fluxes \( \tilde{q}_h \) and \( \tilde{u}_h \) are approximations to \( q_h \) and \( u_h \), respectively, on each mesh face and define how the degrees of freedom in each element are coupled together.

To more succinctly describe the coupling between elements, the following standard notation is adopted. Consider two adjacent elements \( E^+ \) and \( E^- \) which share a face in \( \mathcal{E} \). Let \( n^\pm \) denote the outward unit normals of \( \partial E^\pm \) along the shared face and \((\omega^\pm, v^\pm)\) denote the traces of \((\omega, v) \in V^d_h \times V_h\) from \( E^\pm \) on the shared face. The average \( \{ \cdot \} \) and jump \( [ \cdot ] \) operators on the shared face are then defined as

\[
\{ \omega^\pm \} = \frac{1}{2} ( \omega^+ + \omega^- ), \quad \{ v^\pm \} = \frac{1}{2} ( v^+ + v^- ),
\]

\[
[ \omega^\pm ] = \omega^+ \cdot n^\pm + \omega^- \cdot n^- , \quad [ v^\pm ] = v^+ n^\pm + v^- n^- .
\]

On boundary faces, \((\omega^-, v^-)\) shall refer to the traces of \((\omega, v)\) from the corresponding element touching \( \partial \Omega \).

2.3. The local discontinuous Galerkin method. The choice of numerical flux in (2.5)–(2.6) defines a DG method. Here we focus on the local discontinuous Galerkin method [15], which chooses numerical fluxes \( \tilde{q}_h \) and \( \tilde{u}_h \) according to the general form

\[
\tilde{q}_h = \begin{cases} 
\{ q_h \} + \beta [ q_h ] - \tau_0 [ u_h ] & \text{on any interior face,} \\
q_h - \tau_D ( u_h - g ) n & \text{on any face of } \Gamma_D, \\
h n & \text{on any face of } \Gamma_N,
\end{cases}
\tag{2.7}
\]

and

\[
\tilde{u}_h = \begin{cases} 
\{ u_h \} - \beta \cdot [ u_h ] & \text{on any interior face,} \\
g & \text{on any face of } \Gamma_D, \\
u_h & \text{on any face of } \Gamma_N,
\end{cases}
\tag{2.8}
\]

where \( \beta \) is a (possibly face-dependent) user-defined vector; for example, in a one-sided flux scheme, \( \beta = \pm \frac{1}{2} n \). Here, the numerical flux \( \tilde{q}_h \) includes penalty stabilization terms; \( \tau_0 \geq 0 \) is a penalty parameter associated with interior faces and \( \tau_D > 0 \) is associated with Dirichlet faces (if any). Generally, \( \tau_0 \) must be strictly positive to ensure well-posedness of the discrete problem, but in some cases (e.g., on Cartesian grids with particular choices of \( \beta \)), \( \tau_0 \) can be set equal to zero [12]. If \( \Gamma_D \) is nonempty, \( \tau_D \) must be positive to ensure well-posedness of the final discrete problem. Although arbitrarily small penalty parameters suffice to ensure well-posedness, later we show that choosing sufficiently large values can greatly benefit multigrid performance (see Figure 4.4).

We first particularize (2.5) for the LDG method, which is the weak statement that \( q = \nabla u \). We slightly modify the weak form (2.5) by defining \( q_h \in V^d_h \) in strong-weak...
form,\(^1\) such that

\[
(2.9) \quad \int_E q_h \cdot \omega = \int_E \nabla u_h \cdot \omega + \int_{\partial E} (\tilde{u}_h - u_h) \omega \cdot n
\]

holds for every element \(E \in \mathcal{E}\) and every test function \(\omega \in \{Q_p(E)\}^d\). Upon summing (2.9) over every element of the mesh and using the definition of the numerical flux \(\tilde{u}_h\) in (2.8), we have that, for any \(\omega \in V_h^d\),

\[
(2.10) \quad \int_\Omega q_h \cdot \omega = \sum_{E \in \mathcal{E}} \int_E \nabla u_h \cdot \omega - \int_{\Gamma_0} [u_h] \cdot (\llbracket \omega \rrbracket + \beta \llbracket \omega \rrbracket) + \int_{\Gamma_D} (g - u_h^\tau) \omega^- \cdot n,
\]

where \(\Gamma_0\) denotes the union of all interior faces of \(\mathcal{E}\). Define the following operators:

- Let \(\nabla_h : V_h \rightarrow V_h^d\) be the \textit{broken gradient operator} and \(L : V_h \rightarrow V_h^d\) be the \textit{lifting operator}, such that
  \[
  \int_\Omega (\nabla_h u) \cdot \omega = \sum_{E \in \mathcal{E}} \int_E \nabla u \cdot \omega,
  \]
  \[
  \int_\Omega (Lu) \cdot \omega = -\int_{\Gamma_0} [u] \cdot (\llbracket \omega \rrbracket + \beta \llbracket \omega \rrbracket) - \int_{\Gamma_D} u^- \omega^- \cdot n
  \]
  holds for every \(\omega \in V_h^d\) and each \(u \in V_h\).
- Define \(J_D(g) \in V_h^d\) such that
  \[
  \int_\Omega J_D(g) \cdot \omega = \int_{\Gamma_D} g \omega^- \cdot n
  \]
  holds for every \(\omega \in V_h^d\).

Accordingly, (2.10) is equivalent to the statement that

\[
(2.11) \quad q_h = (\nabla_h + L) u_h + J_D(g) = Gu_h + J_D(g)
\]

where \(G : V_h \rightarrow V_h^d\) is the \textit{discrete gradient operator}, \(G = \nabla_h + L\). The formula (2.11) is the LDG discretization of the statement \(q = \nabla u\), taking into account Dirichlet boundary data.

Next, we particularize (2.6) for the LDG method, which is the weak statement that \(-\nabla \cdot q = f\). Upon summing (2.6) over every mesh element and using the definition of the numerical flux \(\tilde{q}_h\) in (2.7), we have that, for any \(v \in V_h\),

\[
(2.12) \quad \sum_{E \in \mathcal{E}} \int_E q_h \cdot \nabla v - \int_{\Gamma_0} (\llbracket q_h \rrbracket + \beta \llbracket q_h \rrbracket - \tau_0 [u_h]) \cdot [\omega] - \int_{\Gamma_D} (q_h^- \cdot n - \tau_D u_h^-) v^- = \int_\Omega f v + \tau_D \int_{\Gamma_D} g v^- + \int_{\Gamma_N} h v^-.
\]

\(^1\)The \textit{strong-weak form} states that \(q_h\) must satisfy \(\int_E q_h \cdot \omega = \int_E \nabla u_h \cdot \omega + \int_{\partial E} (\tilde{u}_h - u_h) \omega \cdot n\) whereas the \textit{weak form} states that \(q_h\) must satisfy \(\int_E q_h \cdot \omega = -\int_E u_h \nabla \cdot \omega + \int_{\partial E} \tilde{u}_h \omega \cdot n\). The two forms are equivalent whenever the employed quadrature scheme exactly satisfies the identity of integration by parts, which in practice is generally true for quadrilateral, prismatic, simplicial elements, etc., but is generally not true when approximate numerical quadrature schemes are used, e.g., as on implicitly defined curved elements. In the latter situation, to ensure symmetry of the final discrete Laplacian operator, it is necessary to use the strong-weak form to define \(q_h\) and the weak form to define the divergence of \(q_h\) (or vice versa)—see [28] for further remarks.
Additionally, define the following operators:

- Similar to the operator $J_D$ above, let $J_N(h) \in V_h$ be such that
  \[ \int_{\Omega} J_N(h) v = \int_{\Gamma_N} h v^- \]
  for all $v \in V_h$.
- Let $E_0, E_D : V_h \to V_h$ be the operators such that, for each $u \in V_h$,
  \[ \int_{\Omega} E_0(u) v = \int_{\Gamma_0} [u] \cdot [v], \quad \int_{\Omega} E_D(u) v = \int_{\Gamma_D} u^- v^- \]
  holds for every $v \in V_h$. These operators penalize jumps in the discrete solution on interior and Dirichlet boundary faces, respectively.
- Let $a_D(g) \in V_h$ be such that
  \[ \int_{\Omega} a_D(g) v = \int_{\Gamma_D} g v^- \]
  for all $v \in V_h$.

Then, using the fact that $(q_h, \nabla h v) + (q_h, L v) = (q_h, G v)$, (2.12) is equivalent to (2.13)
\[ (q_h, G v) + \tau_0 (E_0 u_h, v) + \tau_D (E_D u_h, v) = (f + J_N(h) + \tau_D a_D(g), v), \]
or, putting aside penalty terms, $G^* q_h = f + J_N(h)$, where $-G^*$ is the discrete divergence operator, the negative adjoint of the discrete gradient operator $G$; this is the LDG discretization of the statement that $-\nabla \cdot q = f$, taking into account Neumann boundary data.

### 2.3.1. Primal formulation.
To obtain the primal formulation of the LDG method, we combine (2.13) with (2.11) to eliminate $q_h$ and arrive at an equation for $u_h$. The primal LDG formulation of (2.1) reads as follows: find $u_h \in V_h$ such that
\[ a(u_h, v) = \ell(v) \]
for all $v \in V_h$, where the bilinear form $a(\cdot, \cdot)$ is given by
\[ a(u_h, v) = (G u_h, G v) + \tau_0 (E_0 u_h, v) + \tau_D (E_D u_h, v) \]
and the linear functional $\ell(\cdot)$ is given by
\[ \ell(v) = (f, v) - (J_D(g), G v) + (J_N(h), v) + \tau_D (a_D(g), v). \]

One may verify that the bilinear form $a(u, v)$ is symmetric. Discretization of the primal form (2.14) with respect to a particular basis of $V_h$ yields a symmetric positive (semi)define linear system of the form
\[ Au_h = \ell, \]
where $A$ is the matrix form of the negative discrete Laplacian operator.

---

2 Throughout the paper we shall frequently use the same symbol to denote (i) elements of spaces such as $V_h$ or operators acting on such elements, and (ii) vectors of coefficients in the chosen basis or matrices acting on such vectors. The distinction should be clear from context. Furthermore, the particular choice of basis used in this work is a nodal basis, using the Gauss–Lobatto tensor product nodes for the polynomial spaces $Q_p(E)$. 
**2.3.2. Flux formulation.** An alternative, but equivalent, characterization of the LDG method is the so-called *flux formulation*, which does not eliminate the auxiliary variable \( q_h \) from the system (2.11) and (2.13) but instead retains it as a primary unknown. The flux formulation of (2.1) then reads as follows: find \((q_h, u_h) \in V^d_h \times V_h\) such that

\[
\begin{align*}
\mathbf{m}(q_h, \omega) - \mathbf{grad}(u_h, \omega) &= \mathbf{j}(\omega), \\
-\mathbf{div}(q_h, v) + \tau(u_h, v) &= k(v),
\end{align*}
\]

for all \((\omega, v) \in V^d_h \times V_h\), where

\[
\begin{align*}
\mathbf{m}(q, \omega) &= (q, \omega), \\
\mathbf{grad}(u, \omega) &= (\mathbf{G}u, \omega), \\
\mathbf{div}(q, v) &= -(q, \mathbf{G}v), \\
\tau(u, v) &= \tau_0(\mathbf{E}_0 u, v) + \tau_D(\mathbf{D} u, v), \\
\mathbf{j}(\omega) &= (\mathbf{J}_D(g), \omega), \\
k(v) &= (f, v) + (\mathbf{J}_N(h), v) + \tau_D(a_D(g), v).
\end{align*}
\]

Discretization of the flux form (2.16) with respect to a particular basis of \( V^d_h \times V_h \) yields a symmetric positive (semi)definite linear system of the form

\[
\begin{bmatrix}
\mathbf{M} & -\mathbf{M}G \\
-\mathbf{M}D & \mathbf{M}T
\end{bmatrix}
\begin{bmatrix}
q_h \\
u_h
\end{bmatrix}
= \begin{bmatrix}
\mathbf{j} \\
k
\end{bmatrix}.
\]

Here, \( \mathbf{M} \) is the block diagonal mass matrix for \( V_h \) or \( V^d_h \) (depending on context), \( \mathbf{G} \) is the matrix form of the discrete gradient operator, \( \mathbf{D} = -\mathbf{M}^{-1}\mathbf{G}^\top \mathbf{M} \) is the matrix form of the discrete divergence operator, and \( T = \tau_0 \mathbf{E}_0 + \tau_D \mathbf{E}_D \) contains the discrete penalty terms. Since \( M^{-1} \) is also block diagonal, we can easily take the Schur complement of \( M \) in (2.17) to obtain a linear system for the unknown vector \( u_h \),

\[
\mathbf{A} u_h = \ell,
\]

where \( \mathbf{A} = \mathbf{M}(-\mathbf{D}G + T) = \mathbf{G}^\top \mathbf{M}G + \mathbf{M}T \) and \( \ell = k - \mathbf{G}^\top \mathbf{j} \).

The reduced linear system (2.18) is equivalent to the discrete primal formulation (2.15). However, as we will demonstrate next, the two formulations have different implications for multigrid methods. In particular, applying standard operator coarsening to the discrete primal formulation results in poor multigrid performance; coarsening the discrete flux formulation (2.17) in both \( q_h \) and \( u_h \) before taking the Schur complement (2.18) results in optimal multigrid performance, and is equivalent to pure geometric multigrid.

**3. Multigrid methods.** We assume here that the reader has some familiarity with multigrid methods; see for example Briggs, Henson, and McCormick [11] for a review of their design and operation. A geometric multigrid method consists of four main ingredients: a mesh hierarchy, an interpolation operator to transfer approximate solutions from a coarse mesh onto a fine mesh, a restriction operator to formulate a coarse mesh correction problem by restricting the residual from the fine mesh, and a smoother/relaxation method. We consider these ingredients separately first, and then combine them into a multigrid V-cycle.

In the following sections we focus our description on \( h \)-multigrid methods, wherein the mesh is coarsened geometrically at each level. However, much of our analysis
carries over to \(p\)-multigrid methods, which hold the mesh fixed and instead coarsen the
polynomial space by reducing \(p\) at each level. We will try to point out the distinctions
between the two methods when the analogues are not immediately obvious, though
we will use the notation \(h\) in our description.

3.1. Mesh hierarchy. In this work, we employ quadtrees (in 2D) and octrees
(in 3D) to define the finest mesh—whether it is uniform, adaptively refined, or used as
the background grid for an implicitly defined mesh (see section 4.3). The tree structure
then naturally defines a hierarchy of nested meshes for use in \(h\)-multigrid, which are
spatially coarsened by a factor of two in each dimension on each level. For adaptively
refined meshes where the cell size is not uniform, we coarsen each element as rapidly
as the tree structure permits (see, e.g., Figure 4.5).

In the context of \(p\)-multigrid methods, a mesh hierarchy is defined by applying
a specific \(p\)-coarsening strategy to the fine mesh. For example, one could coarsen \(p
sequentially \((p, p-1, p-2, \ldots, 1)\), by a factor of two \((p \to p/2 \to p/4 \to \cdots \to 1)\),
or by some user-defined sequence of \(p\)'s. The first method is a common choice when
low-order polynomials are used on the finest mesh, whereas the second method is
better suited to high-order discretizations.

Mesh hierarchies can also be generated by combining coarsening strategies in both
\(h\) and \(p\). For example, a popular choice is to layer \(p\)-multigrid on top of \(h\)-multigrid,
so that \(h\)-multigrid with a low-order polynomial degree is used as the bottom solver
in the \(p\)-multigrid hierarchy.

3.2. Interpolation. The interpolation operator \(I_f\) transfers a piecewise polynomial
function \(u_c \in V_{2h}(\mathcal{E}_c)\) defined on a coarse mesh to a piecewise polynomial function
\(u_f \in V_h(\mathcal{E}_f)\) on a fine mesh. In the nodal basis, we define \(u_f\) as the corresponding
nodal interpolant of \(u_c\), so that the interpolation operator injects the polynomial defined
on the coarse mesh into the fine mesh unmodified. From the \(h\)-multigrid perspective, \(u_f|_{E_f}\),
on the fine element \(E_f\) is simply the polynomial \(u_c|_{E_c}\) restricted to \(E_f\), where \(E_c \supset E_f\) is the corresponding coarse element in the mesh hierarchy. From
the \(p\)-multigrid perspective, the lower-degree polynomial \(u_c\) can be exactly represented
as a higher-degree polynomial by taking the higher-order coefficients of \(u_f\) to be zero.
Thus, when regarded as an operator from \(L^2(\Omega) \to L^2(\Omega)\), \(I_f\) is the identity operator.
The operator is linear and has the property that it preserves constant functions, i.e.,
\(u_c \equiv 1\) is mapped to \(u_f \equiv 1\). This property ensures that, throughout a V-cycle,
the coarse mesh discrete problems preserve the compatibility condition required in
semidefinite problems having solely Neumann boundary conditions.

3.3. Restriction. We define the restriction operator \(R_f^c : V_h(\mathcal{E}_f) \to V_{2h}(\mathcal{E}_c)\) as
the adjoint of the interpolation operator, i.e., such that
\[
(R_f^c u_f, u_c)_{\mathcal{E}_c} = (u_f, I_f^c u_c)_{\mathcal{E}_f}
\]
holds for every \(u_f \in V_h(\mathcal{E}_f)\) and every \(u_c \in V_{2h}(\mathcal{E}_c)\). Abusing notation and letting \(I_f^c\),
\(R_f^c\), \(u_c\), and \(u_f\) also denote matrices and vectors relative to the nodal bases of \(V_h(\mathcal{E}_f)\)
and \(V_{2h}(\mathcal{E}_c)\), (3.1) is equivalent to
\[
(R_f^c u_f)^\top M_c u_c = u_f^\top M_f I_f^c u_c
\]
where \(M_c\) and \(M_f\) are the block-diagonal mass matrices of the coarse and fine meshes,
respectively. Thus, in the nodal basis,
\[
R_f^c = M_c^{-1} (I_f^c)^\top M_f.
\]
Defining the restriction operator in this manner—sometimes referred to as Galerkin projection—results in several notable properties:

- One may interpret $R_f^h u_f$ as “averaging” elemental polynomials of $u_f \in V_h(E_f)$ on the fine mesh to determine a coarsened piecewise-polynomial representation on the coarse mesh. The averaging is performed in a way that locally preserves the mass of $u_f$: indeed, since $I^f_c$ preserves constant functions, we have that $(R_f^h u_f, 1)_{E_c} = (u_f, 1)_{E_f}$ for all $u_f \in V_h(E_f)$.
- The adjoint method can also be viewed as an $L^2$ projection of $u_f \in V_h(E_f)$ onto $V_{2h}(E_c)$. The variational problem arg min$_{u_c \in V_{2h}(E_c)} \|u_c - u_f\|_0^2$ optimizes the functional
  \[
  V_{2h}(E_c) \ni u_c \mapsto (u_c, u_c)_{E_c} - 2(u_f, I_c^f u_c)_{E_f} = u_c^\top M_c u_c - 2 u_f^\top M_f I_c^f u_c
  \]
  whose unique minimum is given by $u_c = M_c^{-1}(I_c^f)^\top M_f u_f$.
- From the preceding property, one can immediately infer that
  \[
  (3.3) \quad R_f^h I^f_c = I_c^n
  \]
  where $I_c^n$ is the identity operator; i.e., interpolating a piecewise polynomial function from a coarse mesh onto a fine mesh and immediately restricting the result shall return the original function. The relation in (3.3) together with (3.2) also provides a method to compute the coarse-mesh mass matrix from the fine-mesh mass matrix:
  \[
  (3.4) \quad M_c = (I_c^f)^\top M_f I_c^f.
  \]

Given a linear operator $A : V_h(E_f) \rightarrow V_h(E_f)$, one can define a coarsened operator $C(A) : V_{2h}(E_c) \rightarrow V_{2h}(E_c)$ in a similar way, by proceeding variationally: we define $C(A)$ such that
  \[
  (C(A)u_c, v_c)_{E_c} = (A I_c^f u_c, I_c^f v_c)_{E_f}
  \]
holds for all $u_c, v_c \in V_{2h}(E_c)$. Viewing $A$ and $C(A)$ as matrix operators, mapping vectors in the nodal bases of $V_h(E_f)$ and $V_{2h}(E_c)$,
  \[
  C(A) = M_c^{-1}(I_c^f)^\top M_f A I_c^f = R_f^h A I_c^f.
  \]
The last form is perhaps more commonly seen or referred to as “RAT” in the multigrid literature [34] (where $R$ is restriction, $A$ is the fine-mesh operator, and $T$ (or $P$) is the interpolation (or prolongation) operator): the essence of the present work is to show that directly applying $RAT$ to the negative discrete Laplacian resulting from the primal formulation of an LDG method results in an inefficient multigrid algorithm and that, instead, applying $RAT$ to the flux formulation, $q = \nabla u$, $-\nabla \cdot q = f$, leads to more efficient multigrid solvers.

### 3.4. Operator coarsening and pure geometric multigrid.

In this section we compare a standard, purely geometric multigrid method to two multigrid schemes based on operator coarsening: (i) applying $RAT$ to the negative discrete Laplacian matrix of the primal formulation (“primal coarsening”) and (ii) applying $RAT$ to the $2 \times 2$ block matrix of the flux formulation (“flux coarsening”). By a pure geometric method, we mean one in which each level of the hierarchy is explicitly meshed and
the LDG formulation is canonically applied to each level, with the above restriction and interpolation operators transferring residual and correction vectors (in a V-cycle) between levels. Our motivation here concerns an \( h \)-multigrid method; however, much of the following discussion has direct analogy with \( p \)-multigrid methods. With this in mind, we discuss interaction between a pair of levels: suppose \( \mathcal{E}_f \) is the mesh of a fine level and \( \mathcal{E}_c \) is the mesh of the next-coarsest level.

### 3.4.1. Primal coarsening

Recall the primal form of the negative discrete Laplacian operator of an LDG method: as a matrix (mapping the nodal basis of \( V_h(\mathcal{E}_f) \) into the nodal basis of \( V_h(\mathcal{E}_f) \)),

\[
A = -DG + \tau_0 E_0 + \tau_D E_D,
\]

where \( G = \nabla_h + L \) is the discrete gradient operator and \( D = -M^{-1}G^TM \) is the discrete divergence operator. To discuss the application of \( RAT \) to \( A \) and how it relates to a geometric multigrid implementation, we consider the individual terms making up \( A \). In the following, operators with subscripts \( f, h \) and \( c, 2h \) denote those constructed with LDG on the fine mesh and coarse mesh, respectively.

- First, the broken gradient operator satisfies the \( RAT \) property, i.e., \( \mathcal{C}(\nabla_h) = \nabla_{2h} \). Computing the piecewise gradient on a coarse mesh and interpolating the result to the fine mesh is the same as computing the piecewise gradient of the interpolant, i.e., \( I_f^T \nabla_{2h} u_c = \nabla_h I_c^T u_c \) for all \( u_c \in V_{2h}(\mathcal{E}_c) \); consequently, \( \mathcal{C}(\nabla_h) = R^T f \nabla_h I_c^T = R^T f I_f \nabla_{2h} = \nabla_{2h} \) by (3.3).

- The lifting operator also satisfies the \( RAT \) property, i.e., \( \mathcal{C}(L_f) = L_c \). This is perhaps not immediately obvious, since source terms on a coarse mesh face will lift into the corresponding large coarse element, whereas the corresponding source terms on the fine mesh faces lift only into the smaller elements touching that face; however, the restriction of the result on the set of smaller elements agrees with the result of \( L_c \). To see this, apply the variational formulation of \( \mathcal{C}(\cdot) \) to observe that

\[
(L_f I_c^T u_c, I_c^T v_c)_{\mathcal{E}_c} = (L_f I_c^T u_c, I_c^T v_c)_{\mathcal{E}_f} = \int_{\Gamma_{0,f}} \|I_f^T u_c\| \cdot (\|I_c^T v_c\| + \beta \|I_c^T v_c\|) - \int_{\Gamma_{D,f}} (I_f^T u_c)^- (I_f^T v_c)^- \cdot n = -\int_{\Gamma_{0,c}} \|u_c\| \cdot (\|v_c\| + \beta \|v_c\|) - \int_{\Gamma_{D,c}} u_c^- v_c^- \cdot n = (L_c u_c, v_c)_{\mathcal{E}_c}
\]

holds for all \( u_c \in V_{2h}(\mathcal{E}_c) \) and \( v_c \in V_{2h}^d(\mathcal{E}_c) \). Here, \( \Gamma_{0,f} \) and \( \Gamma_{0,c} \) denote the union of interior faces of the fine and coarse meshes, respectively, and similarly for \( \Gamma_{D,f} \) and \( \Gamma_{D,c} \). The third equality holds because the interpolation operator introduces no nonzero jumps on the set of new fine mesh faces, i.e., on \( \Gamma_{0,f} \setminus \Gamma_{0,c} \) and \( \Gamma_{D,f} \setminus \Gamma_{D,c} \). (The preceding assumes that fine mesh faces inherit the same \( \beta \) value as coarse mesh faces; in particular, this is true for the one-sided LDG scheme in which \( \beta = \pm \frac{1}{2}n \).)

- It immediately follows from the preceding two properties that \( \mathcal{C}(G_f) = G_c \).
Moreover,
\[
C(D_f) = R_f^c D_f I_c^f = (M_c^{-1}(I_c^f)^\top M_f)(-M_f^{-1}G_f^\top M_f)I_c^f
\]
\[
= -M_c^{-1}((I_c^f)^\top G_f^\top M_f I_c^f M_c^{-1})M_c = -M_c^{-1}(C(G_f))^\top M_c
\]
\[
= -M_c^{-1}G_c^\top M_c = D_c.
\]

- It is straightforward to show that the penalty terms also satisfy the RAT property, i.e., \(C(E_{0,f}) = E_{0,c}\) and \(C(E_{D,f}) = E_{D,c}\). As in the case of the lifting operator, this property derives from the fact the interpolation operator does not introduce jumps on fine mesh faces that do not overlap with coarse mesh faces.

Despite these consistencies, the negative discrete Laplacian does not satisfy the RAT property—the application of RAT to the fine-mesh negative discrete Laplacian \(A_f\) does not yield the coarse-mesh operator \(A_c\) obtained from pure geometric multigrid. Using the properties derived above,
\[
A_c = -D_c G_c + \tau_0 E_{0,c} + \tau_D E_{D,c}
\]
\[
= -C(D_f)C(G_f) + \tau_0 C(E_{0,f}) + \tau_D C(E_{D,f})
\]
which differs from the direct coarsening of \(A_f\),
\[
C(A_f) = -C(D_f G_f) + \tau_0 C(E_{0,f}) + \tau_D C(E_{D,f}) \neq A_c,
\]

since in general \(C(D_f G_f) \neq C(D_f)C(G_f)\). Informally, \(C(D_f G_f)u_c\) interpolates a function \(u_c \in V_{2h}(\mathcal{E}_c)\) onto the fine mesh \(\mathcal{E}_f\), computes the gradient as a function in \(V_h^d(\mathcal{E}_f)\), computes the divergence as a function in \(V_h(\mathcal{E}_f)\), and projects the result (via restriction) back to the coarse mesh \(\mathcal{E}_c\). On the other hand, \(C(D_f)C(G_f)u_c\) projects the computed fine-mesh gradient onto the coarse mesh (via restriction) and then immediately interpolates the result in order to compute the discrete divergence on the fine mesh, before projecting the final result back to the coarse mesh. That is,
\[
C(D_f)C(G_f) = C(D_f I_c^f R_f^c G_f) \neq C(D_f G_f),
\]
since \(I_c^f R_f^c \neq I\).

### 3.4.2. Flux coarsening

The true coarse operator \(A_c\) obtained from pure geometric multigrid may be viewed as applying RAT to the equations \(q = \nabla u\) and \(-\nabla \cdot q = f\) separately. The flux formulation of LDG, (2.17), naturally displays this coarsening strategy. To show this, note that we can write the flux formulation—which operates in the mass basis—in the nodal basis by premultiplying (2.17) by the inverse mass matrix to obtain
\[
(3.5) \quad \begin{bmatrix} I & -G \\ -D & T \end{bmatrix} \begin{bmatrix} q_h \\ u_h \end{bmatrix} = \begin{bmatrix} M_{-1, f} \end{bmatrix}.
\]

Applying RAT in a block fashion to the nodal flux formulation (3.5) then yields the discrete operator
\[
(3.6) \quad \begin{bmatrix} R_f^c & 0 \\ 0 & R_f^c \end{bmatrix} \begin{bmatrix} I & -G_f \\ -D_f & T_f \end{bmatrix} \begin{bmatrix} I_c^f \\ 0 \end{bmatrix} = \begin{bmatrix} I & -G_c \\ -D_c & T_c \end{bmatrix}.
\]
Taking the Schur complement of (3.6), we obtain

\[ A_c = -D_c G_c + T_c \]

\[ = -C(D_f)C(G_f) + \tau_0 C(E_{0,f}) + \tau_D C(E_{D,f}), \]

which is exactly the coarse operator from pure geometric multigrid. Thus, applying operator coarsening to the flux formulation of LDG, which is equivalent to separately coarsening the equations \( q = \nabla u \) and \( -\nabla \cdot q = f \), is the same as pure geometric multigrid.

Figure 3.1 depicts the three types of coarsening that can be performed, given a hierarchy of meshes. In the left column, pure geometric multigrid defines the coarse operators directly from the coarse meshes; in the center column, primal coarsening applies \( RAT \) to the discrete Laplacian operator; and in the right column, flux coarsening applies \( RAT \) separately to the discrete divergence and gradient operators. In the above, we have shown the equivalence of the left and right columns. An implementation of constructing the operator hierarchy using flux coarsening is outlined in Algorithm 3.1.

### 3.4.3. Benefits of operator coarsening

It can be useful to define coarse operators directly from fine operators (e.g., by using \( RAT \)) rather than via discretizations computed directly from coarse meshes. Since coarse mass matrices can be computed automatically according to (3.4), quadrature schemes do not need to be computed for coarse elements—instead, fine-mesh quadrature rules are coarsened automatically via (3.4). Similarly, coarse lifting matrices are not explicitly needed since their contribution to the discrete gradient is automatically computed via \( G_c = C(G_f) \), and so quadrature rules for coarse faces also do not need to be defined. For implicitly defined meshes, such as the ones shown in section 4.3, computing coarse quadrature rules can be computationally intricate or taxing; the fact that operator coarsening obviates
the need for this is a substantial benefit. Additionally, operator coarsening can be efficiently implemented using basic linear algebra operations (e.g., block-sparse matrix multiplication) for which highly optimized and parallelized libraries exist; in contrast, computing discretizations directly from coarse meshes relies heavily on the efficiency of one’s own code.

3.4.4. Relation to other DG methods. Although we have focused on the LDG method in our discussion, we expect that other DG methods may require similar care in coarsening fine-grid operators such that they are consistent with a pure geometric multigrid method. For instance, methods for which the numerical flux $\hat{q}_h$ depends on the discrete gradient of $u_h$—such as the BR1 [6] or Brezzi [10] methods—may need similar treatment, as the contribution from the lifting operator $L$ must be coarsened separately.

Other methods, such as the symmetric interior penalty (SIP) method, do not require the discrete divergence and discrete gradient operators to be coarsened separately. To demonstrate this for SIP, start with its corresponding bilinear form for a pure Neumann problem: find $u \in V_h$ such that
$$a(u, v) = l(v) \quad \text{for all } v \in V_h,$$
where
$$a(u, v) = (\nabla_h u, \nabla_h v) - \int_{\Gamma_0} (\{\nabla_h u\} \cdot \{v\} + [u] \cdot \{\nabla_h v\}) + \tau \int_{\Gamma_0} [u] \cdot [v]$$
and
$$l(v) = (f, v) + \int_{\Gamma_N} h \, v^-. $$
Now consider a pure geometric multigrid method. Let $u_c, v_c \in V_{2h}(E_c)$. Then
$$a_c(u_c, v_c) = (\nabla_{2h} u_c, \nabla_{2h} v_c) - \int_{\Gamma_{0,c}} (\{\nabla_{2h} u_c\} \cdot \{v_c\} + [u_c] \cdot \{\nabla_{2h} v_c\})$$
$$+ \tau \int_{\Gamma_{0,c}} [u_c] \cdot [v_c]$$
$$= (\nabla_h I_c u_c, \nabla_h I_c v_c) - \int_{\Gamma_{0,f}} (\{\nabla_h I_c u_c\} \cdot \{I_c v_c\} + [I_c u_c] \cdot \{\nabla_h I_c v_c\})$$
$$+ \tau \int_{\Gamma_{0,f}} [I_c u_c] \cdot [I_c v_c]$$
$$= a_f(I_c u_c, I_c v_c),$$
where equality holds between the first and second lines because $I_c$ does not introduce nonzero jumps on new mesh faces. Therefore, as matrices (mapping vectors in the nodal basis to vectors in the nodal basis),
$$u_c^\top M_c A_c v_c = (I_c u_c)^\top M_f A_f (I_c v_c).$$
Assuming the quadratic form is nondegenerate (which is true since $A_c$ and $A_f$ are symmetric positive definite, ignoring the trivial kernel), this implies
$$A_c = M_c^{-1}(I_c)^\top M_f A_f I_c = R_f A_f I_c.$$
This is $RAT$ applied to $A_f$, and so applying pure geometric multigrid to SIP is the same as recursively applying standard (primal) operator coarsening to $A_f$. 

13
Algorithm 3.1 Construction of coarse operators, Build($M_f, G_f, T_f$)

<table>
<thead>
<tr>
<th>Input:</th>
<th>List of coarse operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>Output:</td>
<td></td>
</tr>
<tr>
<td>$A := {}$</td>
<td></td>
</tr>
<tr>
<td>$A_f := G_f^T M_f G_f + T_f$</td>
<td></td>
</tr>
<tr>
<td>if $\mathcal{E}_f$ is not the coarsest mesh then</td>
<td></td>
</tr>
<tr>
<td>$M_c := (I_f^e)^T M_f I_f^e$</td>
<td></td>
</tr>
<tr>
<td>$G_c := M_f^{-1} (I_f^e)^T M_f G_f I_f^e$</td>
<td></td>
</tr>
<tr>
<td>$T_c := (I_f^e)^T T_f I_f^e$</td>
<td></td>
</tr>
<tr>
<td>$A := \text{Build}(M_c, G_c, T_c)$</td>
<td></td>
</tr>
<tr>
<td>return ${A_f, A}$</td>
<td></td>
</tr>
</tbody>
</table>

Algorithm 3.2 Multigrid V-cycle $V(x_f, b_f)$ on mesh $\mathcal{E}_f$ with $\nu$ pre- and post-smoothing steps

1: if $\mathcal{E}_f$ is the coarsest mesh then
2: Solve $A_f x_f = b_f$ directly
3: else
4: Relax forward $\nu$ times
5: $r_c := (I_f^e)^T (A_f x_f - b_f)$
6: $x_c := V(0, r_c)$
7: $x_f := x_f - I_f^e x_c$
8: Relax backward $\nu$ times
9: return $x_f$

3.5. Multigrid preconditioned conjugate gradient. Recall that a geometric multigrid method utilizes a combination of relaxation/smoothing together with interpolated approximate solutions of coarse-grained problems. In the present case, the coarse-grained problem solves for the correction in a residual equation for the same elliptic problem except on a coarser mesh. In particular, it is important to note that the LDG methods developed in section 2.3 operate in the mass basis, i.e., instead of solving $-\Delta_h u = f$, where $\Delta_h$ is the discrete Laplacian, one instead solves $-M \Delta_h u = M f$, as the latter system is symmetric positive (semi)definite. Thus, to appropriately define the coarse mesh problem, one may: (i) calculate the residual of the fine mesh linear system $A_f x_f = b_f$, (ii) multiply the residual by $M_f^{-1}$ to correctly determine the residual as a piecewise polynomial function, (iii) restrict the residual to the coarse mesh, and then (iv) multiply this residual by $M_c$ of the coarse mesh. Thus, the coarse mesh problem consists of (approximately) solving for $u_c$ such that

$$A_c x_c = M_c (R_f^e [M_f^{-1} (A_f x_f - b_f)]),$$

which, according to the derived restriction operator (3.2), conveniently simplifies to

$$A_c x_c = (I_f^e)^T (A_f x_f - b_f),$$

and so it is unnecessary to multiply by mass matrices in the implementation of the multigrid method; instead, one can simply apply the transpose of the interpolation matrix. With this consideration in mind, the design of a multigrid V-cycle is relatively straightforward and is outlined in Algorithm 3.2.

The V-cycle is designed to preserve the symmetric positive (semi)definite property of the discrete problem, making it suitable for preconditioning the conjugate gradient method. To that end, the relaxation sweeps are performed in a symmetric fashion; the first set of relaxation sweeps uses a given ordering of the unknowns, and the second set uses the reverse of that ordering (referred to as forward and backward relaxation in Algorithm 3.2). A multigrid preconditioned conjugate gradient method [33] (MGPCG) combines the advantages of both solvers: the multigrid preconditioner is effective in the interior of the domain where the elliptic behavior of the matrix dominates, while the conjugate gradient method effectively treats the remaining eigenmodes associated with imposing boundary conditions (and in the case of multi-phase elliptic interface problems, jump conditions on internal interfaces) [32, 28]. We use a single multigrid V-cycle as a precon convergence rate.
4. Numerical results. In this section, we present numerical experiments to assess the efficacy of flux coarsening for LDG discretizations of elliptic PDEs. As the smoother/relaxation method, we use a weighted block Gauss–Seidel smoother with weight $\omega = 0.8$ and $\nu = 3$ pre- and post-smoothing steps in the V-cycle. We set the interior penalty parameter to $\tau_0 = 0.1$, as we found smaller $\tau_0$ to result in better multigrid performance. We measure multigrid performance via the average convergence factor

$$\rho = \exp \left( \frac{1}{N} \log \frac{\|e_N\|_2}{\|e_0\|_2} \right),$$

where $N$ is the number of iterations required to reduce the relative error by a factor of $10^{-10}$ and $e_i$ is the error at iteration $i$ of either standalone multigrid (i.e., $i$ many V-cycles) or MGPCG (i.e., the $i^{th}$ iteration of CG preconditioned by a single V-cycle). In effect, $\rho$ measures the average slope of $e_i$ on a log-linear graph. Convergence is measured using a right-hand side of $f = 0$ with a random nonzero initial guess for $u$.

4.1. Uniform Cartesian grids. We start by solving (2.1) with homogeneous Neumann boundary conditions on the domain $\Omega = [0,1]^d$ using a uniform Cartesian grid of size $n \times n$ (for $d = 2$) or $n \times n \times n$ (for $d = 3$) with grid spacing $h = 1/n$. We build an $h$–multigrid hierarchy based on uniform grid refinement by applying both primal and flux coarsening to the discretized LDG system, and solve using both standalone V-cycles and MGPCG.

Figures 4.1 and 4.2 show the average convergence factor versus $n$ for polynomial orders $1 \leq p \leq 5$ in 2D and 3D, respectively. In both cases, the multigrid scheme built using flux coarsening exhibits $h$-independent convergence factors of $\rho \approx 0.1$, whereas the scheme based on primal coarsening exhibits poor performance that degrades as $h \to 0$.

Similar results hold for $p$–multigrid on uniform Cartesian grids. We generate a $p$–multigrid hierarchy by successively halving the polynomial order (i.e., $p \to p/2 \to p/4 \to \cdots \to 1$) and applying both primal and flux coarsening to the discretized LDG system. Figure 4.3 shows convergence factor versus $p$ for grid sizes $4 \leq n \leq 512$ in 2D and 3D. Again, convergence appears to be independent of $p$ (at least up to $p = 8$) when flux coarsening is used with $p$–multigrid, whereas performance degrades with increasing $p$ for primal coarsening.

In our tests, the imposition or combination of Dirichlet, Neumann, or periodic boundary conditions does not affect the conclusions made in this work. However, for problems with Dirichlet boundary conditions the choice of Dirichlet penalty parameter $\tau_D$ can impact multigrid efficiency. Informally, Dirichlet boundary conditions are enforced in a DG method only weakly and the smoothing/relaxation method of a V-cycle can only effectively enforce the boundary condition if the associated penalty parameter is sufficiently strong. Figure 4.4 shows a study of the impact that $\tau_D$ has on multigrid performance for a homogeneous Dirichlet problem on a uniform $n \times n$ mesh. For $\tau_D \leq 100$, the average convergence factor $\rho$ degrades as $n$ increases; a good choice in this case appears to be $\tau_D \approx 1000$. The proper choice of $\tau_D$ is application-dependent.

---

3As is typical in multigrid methods, increasing the number of pre- and post-smoothing steps can increase the convergence efficiency, i.e., decrease $\rho$ as measured by (4.1); however, doing so comes at the cost of a more expensive V-cycle and therefore may be ineffective. We observed that $\nu = 3$ gave the best computational efficiency in our numerical experiments in terms of reducing the error by a fixed factor.
Fig. 4.1: $h$–multigrid convergence factors for primal coarsening ($A_c = \mathcal{C}(A_f)$) and flux coarsening ($A_c = -\mathcal{C}(D_f)\mathcal{C}(G_f) + \mathcal{C}(T_f)$) applied to the LDG discretization of Poisson’s equation on a uniform $n \times n$ Cartesian grid as $h \to 0$. The top row results are computed using V-cycles whereas the bottom row results are computed using MGPCG. The plot markers indicate different polynomial orders: ■, ○, ▲, ×, and ● denote $p = 1, 2, 3, 4,$ and 5, respectively.

Fig. 4.2: $h$–multigrid convergence factors for primal coarsening ($A_c = \mathcal{C}(A_f)$) and flux coarsening ($A_c = -\mathcal{C}(D_f)\mathcal{C}(G_f) + \mathcal{C}(T_f)$) applied to the LDG discretization of Poisson’s equation on a uniform $n \times n \times n$ Cartesian grid as $h \to 0$. The top row results are computed using V-cycles whereas bottom row results are computed using MGPCG. The plot markers indicate different polynomial orders: ■, ○, ▲, ×, and ● denote $p = 1, 2, 3, 4,$ and 5, respectively. (Omitted data points correspond to simulations whose memory requirements approximately exceed 120 GB.)
Fig. 4.3: $p$–multigrid convergence factors for primal coarsening ($A_c = C(A_f)$) and flux coarsening ($A_c = -C(D_f)C(G_f) + C(T_f)$) applied to the LDG discretization of Poisson’s equation on uniform grids. The $p$–multigrid hierarchy is generated by successively halving the polynomial order (i.e., $p \rightarrow p/2 \rightarrow p/4 \rightarrow \cdots \rightarrow 1$). The plot markers indicate different grid sizes: ■, ●, ▲, ×, ♦, ▼, +, and ○ denote $n = 4, 8, 16, 32, 64, 128, 256,$ and $512$, respectively. (Omitted data points correspond to simulations whose memory requirements approximately exceed 120 GB.)

Fig. 4.4: Effect of Dirichlet penalty parameter $\tau_D$ on multigrid performance for an $n \times n$ mesh with $p = 2$.

and concerns not only multigrid performance but also discretization accuracy and effects of penalty stabilization on the conditioning of the linear systems.

4.2. Adaptive mesh refinement. Next, we solve the Neumann problem (2.1) on an adaptively refined Cartesian mesh, where refinement is performed according to some prescribed spatially-varying function. We implement adaptivity using a quadtree (in 2D) or octree (in 3D), which naturally defines a geometric hierarchy of meshes. An example hierarchy is shown in Figure 4.5, where coarsening is performed for each element as rapidly as possible. Note that some elements at various levels of the hierarchy have four neighbors on a single side, so that large elements and small elements may share part of a face.

Figure 4.6 shows the average convergence factor versus $n_{\text{eff}}$ for polynomial orders $1 \leq p \leq 5$, where $n_{\text{eff}}$ is the effective grid size, i.e., the grid size corresponding to the smallest element in the mesh. In both 2D and 3D, the multigrid method based on primal coarsening exhibits performance that degrades as $n_{\text{eff}} \rightarrow \infty$, whereas the
method based on flux coarsening yields good performance independent of $n_{\text{eff}}$ for all $p$ considered.

4.3. Implicitly defined meshes and elliptic interface problems. Our last two examples are designed to exemplify the benefits of operator coarsening by considering cases in which a pure geometric multigrid method would be intricate or difficult to implement, such as on nontrivial domains with complex geometry or for elliptic interface problems in which the interface has extreme geometry. The first example consists of a curved domain containing holes and thin pieces, and the second example is a multi-phase elliptic interface problem with small circles, filaments, and cusps in
Fig. 4.7: A single-phase test case applied to a curved domain using implicitly defined meshes. Depicted is the implied $h$–multigrid hierarchy (elements are randomly colored). The coarse meshes are not explicitly constructed in our operator coarsening approach; instead, the discrete gradient and penalty operators and mass matrices are constructed top-down at each level (see Algorithm 3.1). In particular, note that the bottom level of the hierarchy consists of a mesh with a single element containing two holes.

the interface geometry. In both cases, we make use of a recently developed framework for computing high-order accurate multi-phase multi-physics using implicitly defined meshes \[28, 29\]. The framework shares some aspects with cut-cell techniques wherein a level set function defining the domain geometry or internal interfaces is used to cut through the cells of a background quadtree or octree; tiny cut cells are then merged with neighboring cells to create a mesh in which the shapes of interfacial elements are defined implicitly by the level set function. Quadrature rules for curved elements and nontrivial mesh faces are then computed using high-order accurate schemes for computing integrals on implicitly defined domains restricted to hyperrectangles \[27\]; these quadrature schemes are then used in the DG weak formulation, e.g., for computing mass matrices and the lifting operator $L$ on the finest-level mesh.

In both examples, we consider an elliptic PDE problem with Dirichlet boundary conditions. The Dirichlet penalty parameter is chosen to scale inversely with $h$, the typical element size on the finest mesh, such that $\tau_D = 100/h$; the value 100 was determined empirically as being approximately the smallest possible while giving good multigrid performance. To measure the convergence rate, we apply the MGPCG method to a homogeneous problem with random nonzero initial condition and measure the average convergence rate using (4.1).

The first example of a single-phase Poisson problem on a curved domain is illustrated in Figure 4.7 and consists of a figure-eight domain with two holes surrounded by thin segments. It is important to note that the illustrated mesh hierarchy is implicitly formed by our operator coarsening scheme in Algorithm 3.1 and it is only the finest-level mesh which is built. On coarser levels, the elements are agglomerated according to the coarsening of the background quadtree; in particular, we note that the bottom level consists of a single element containing two holes. Using the operator coarsening strategy, there is no need to compute quadrature rules for the coarse levels of the mesh hierarchy—the quadrature rules from the fine mesh are effectively coarsened automatically. The results for solving the Dirichlet problem (2.1) using flux coarsening and $h$–multigrid on the curved domain of Figure 4.7 are shown in Figure 4.9 (left); we observe good multigrid convergence factors of $\rho \approx 0.1–0.2$, independent of grid size, for $1 \leq p \leq 5$. 19
Fig. 4.8: A test case involving a multi-phase elliptic interface problem. Depicted is the implied $h$–multigrid hierarchy, wherein the elements are randomly colored a shade of green or blue for phase one or two, respectively. Similar to Figure 4.7, only the finest mesh is explicitly constructed; on coarser levels, the discrete gradient and penalty operators and mass matrices are constructed top-down at each level (see Algorithm 3.1). Note that the bottom level of the hierarchy consists of a mesh of just two elements—the blue element has one component with five holes, whereas the green element consists of six connected components. See the discussion following (4.2) for a description of the noted features A, B, C, D$_1$, and D$_2$.

The second example considers a two-phase elliptic interface problem in a rectangular domain illustrated in Figure 4.8. The corresponding PDE consists of solving

$$\begin{align*}
-\nabla \cdot (\mu_1 \nabla u) &= f \quad \text{in } \Omega_1, \\
[u] &= g_I \quad \text{on } \Gamma, \\
[u] &= g_\partial \quad \text{on } \partial \Omega, \\
[\mu \nabla u \cdot n] &= h_I \quad \text{on } \Gamma,
\end{align*}$$

(4.2) where $\Gamma$ is the interface between phases $\Omega_1$ (green region, with ellipticity coefficient $\mu_1 = 1$) and $\Omega_2$ (blue region, with $\mu_2 = 4$), and $[\cdot]$ denotes the interfacial jump in the indicated quantity. In this test case, the geometry of the interface has been designed to be challenging—the crescent shape is long and thin; there are three isolated, small circles; and the star shape has sharp cusp-like corners. Once more we note that the illustrated hierarchy in Figure 4.8 is implicitly formed by the operator coarsening.
Fig. 4.9: $h$–multigrid convergence factors for flux coarsening and MGPCG applied to the LDG discretization of (left) the single-phase Poisson problem on the curved domain shown in Figure 4.7 and (right) the multi-phase elliptic interface test problem in (4.2) on the domain shown in Figure 4.8. The grid size $n$ is the number of cells of the background Cartesian grid required to cover the longest extent of the domain. The plot markers indicate different polynomial orders: ■, ●, ▲, ×, and ● denote $p = 1, 2, 3, 4$, and 5, respectively.

strategy, and only the finest-level mesh is actually built. However, we note that the agglomeration strategy for this multi-phase problem is slightly different from the previous test examples—here, elements are only agglomerated with elements belonging to the same phase. Thus, the interface remains sharp throughout all levels of the hierarchy, and this can dramatically improve the performance of multigrid methods for elliptic interface problems, especially when using high-order accurate techniques [28]. Owing to this agglomeration strategy, coarse mesh levels can have intricate element shapes. Some example features are noted in Figure 4.8—A indicates a tiny green-phase element surrounded by a large blue-phase element; B indicates two sliver elements; C indicates an element whose cusp-like corners would make it rather difficult to apply a black-box quadrature scheme if one were to explicitly build the coarse-level mesh; and D1,2 show two green-phase elements, each with multiple connected components (three for D1 and four for D2), which is perhaps rather unusual for a finite element method. Another aspect which motivated the present work on operator coarsening is that it would be nearly impossible to directly apply the cell merging algorithms underlying implicitly defined meshes (see [28]) to these coarse levels. Although elements with extreme shapes like these (especially tiny elements next to large elements) can traditionally be of concern for numerical discretization of PDEs, according to our tests they pose no problem when present in the coarse levels of a multigrid solver. Results for solving the homogeneous version of (4.2) ($f$, $g_\partial$, $g_\Gamma$, and $h_\Gamma$ all zero) with a random nonzero initial guess using flux operator coarsening are shown in Figure 4.9 (right). For an elliptic interface problem with challenging interface geometry such as this, convergence factors of $\rho \approx 0.2–0.4$ are reasonable.

5. Concluding remarks. We have presented an $hp$–multigrid method for DG discretizations of elliptic problems that is based on coarsening the discrete gradient and divergence operators from the flux formulation. We have shown that coarsening fine-grid operators in this way results in a method that is equivalent to pure geometric multigrid, but avoids the need to compute bothersome quantities associated with coarse meshes, such as quadrature rules. Whereas traditional Galerkin operator coarsening applied to the primal formulation exhibits poor multigrid performance, operator coarsening applied to the flux formulation performs well independent of both mesh size $h$ and polynomial order $p$, as we demonstrated for uniform Cartesian
grids, adaptively refined meshes, and implicitly defined meshes on complex geometries. Applying the ideas presented here to problems involving more general unstructured meshes is currently under investigation.

Though most of our analysis has focused on the LDG method, we believe that the essential observation applies to other forms of DG discretization of elliptic problems, particularly those in which lifting operators enter the numerical flux for $q$. A more thorough analysis for other DG methods and more general choices of numerical fluxes would be required to determine whether the multigrid method described here extends to other methods, such as CDG or HDG. The idea of coarsening the divergence and gradient operators separately may also be useful for AMG methods, which currently treat the discrete Laplacian operator in its entirety as a black box. Indeed, black-box AMG algorithms applied to LDG discretizations appear to struggle [24]; perhaps applying AMG separately to the divergence and gradient operators in the flux formulation may yield better results.

REFERENCES


